

Last week

- Grading labs & test 2
- Do have lab; if you are finished, let me know
- Office hours
- "Final" goes out on 18th
- Extended hours

Register Allocation

Registers are faster than cache

mov [rsp + 16], rax

Registers are faster than memory

mov rcx, rax

~2-3 orders of magnitude

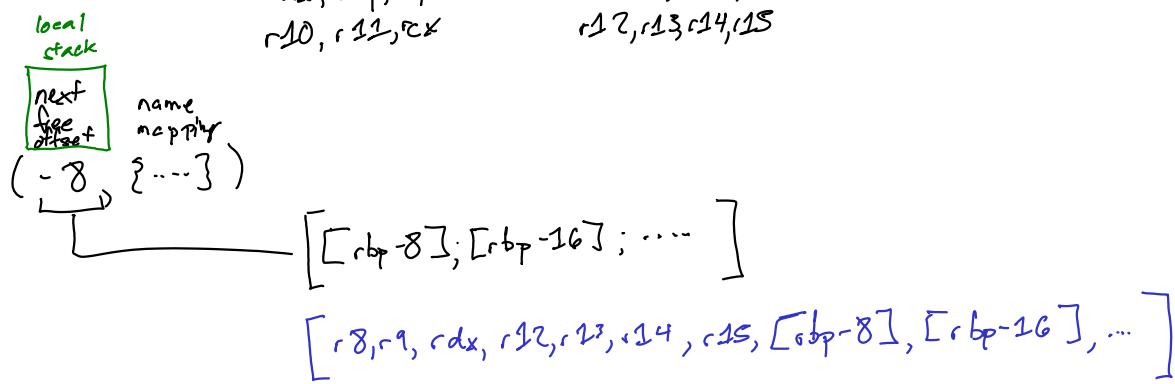
Algorithms

* Greedy

* Linear scanning

* Graph coloring

C-Greedy



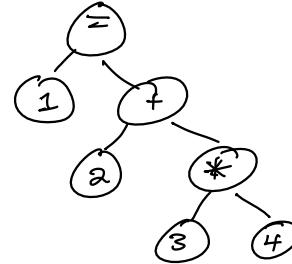
$$\text{Example: } 1 = 2 + 3 * 4$$

Suppose: Storage reg: $r8, r9, rdx$ only

$[r8, r9, rdx, [rbp-8], \dots]$

```

mov rax, 2
mov r8, rax
mov rax, 4
mov r9, rax
mov rax, 6
mov rdx, rax
mov rax, 8
mov [rbp-8], rax
mov rax, rdx
imul rax, [rbp-8]
mov rdx, rax
;
```



Problem case

```

let x = 1 in
let y = 2 in
let z = 3 in
let answer =
    ; (scratches)
in
x + y + z + answer
;
```

Option:

- Pre-analyze code to determine what should go where
- Recognize that we need more space while generating

Linear scanning

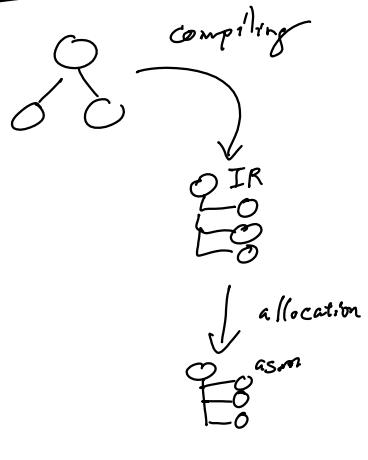
$$1 = 2 \pm 3 * 4$$

fricky

Most compilers do not directly generate assembly. Most compilers use an intermediate representation (IR).

```

mov    ans, 2
mov    loc(1), ans
mov    ans, 4
mov    loc(2), ans
mov    ans, 6
mov    loc(3), ans
mov    ans, 8
mov    loc(4), ans
mov    ans, loc(3)
imul   ans, loc(4)
;
```



Graph Coloring

"interference"

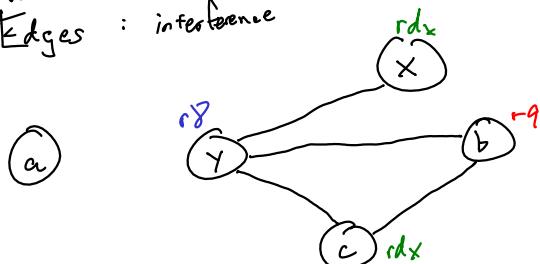
```

let x = 1   in {x}
let y = 2*x in {x,y}
let a = x   in {x,y}
let b = y*2 in {y}
let c = 3   in {y,b}
y+b+c      in {y,b,c}

```

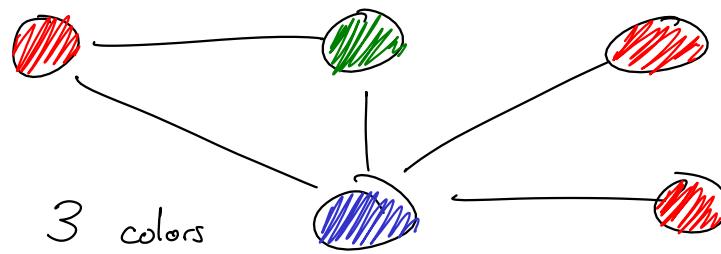
Offices; abstract storage locations

Edges : interference



Color graph

Color represents a concrete storage location



C_{coal} : color vertices s.t. no two neighbors have
 - same color, minimizing colors

Generally : NP-complete
 Colors ≥ 3

Good polynomial approximations

gen code $\sim 125^\circ$ faster than linear scanning