

# Tail Call Optimization

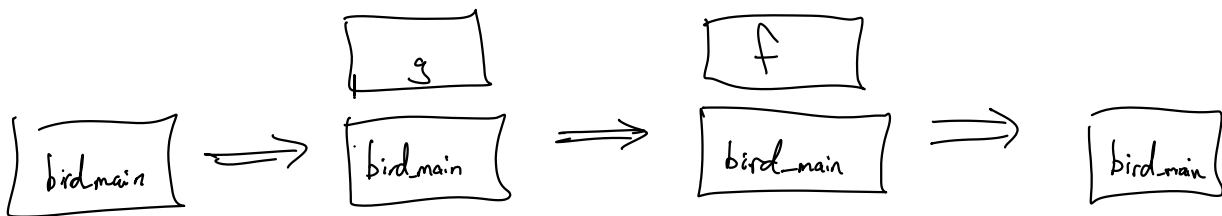
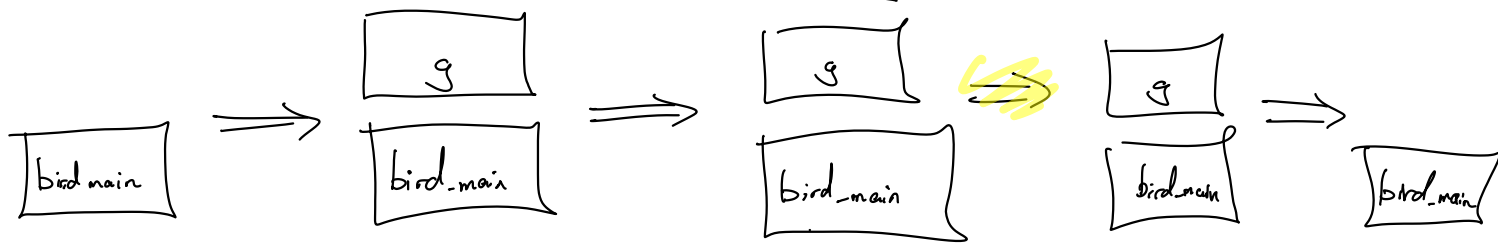
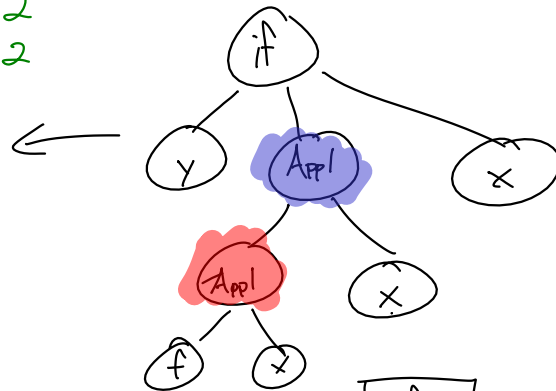
What is? When performing a call, we can remove a stack frame before creating a new one.

Why? Loops = Recursion, recursion eats stack memory  
 $\Rightarrow$  Limit on size loop

Save stack memory using TCO; run programs that can't run otherwise

```
def f a b
  a + b
end
def g x y = (f x) x
end
g 5 true
```

f  $\mapsto$  2  
 g  $\mapsto$  2

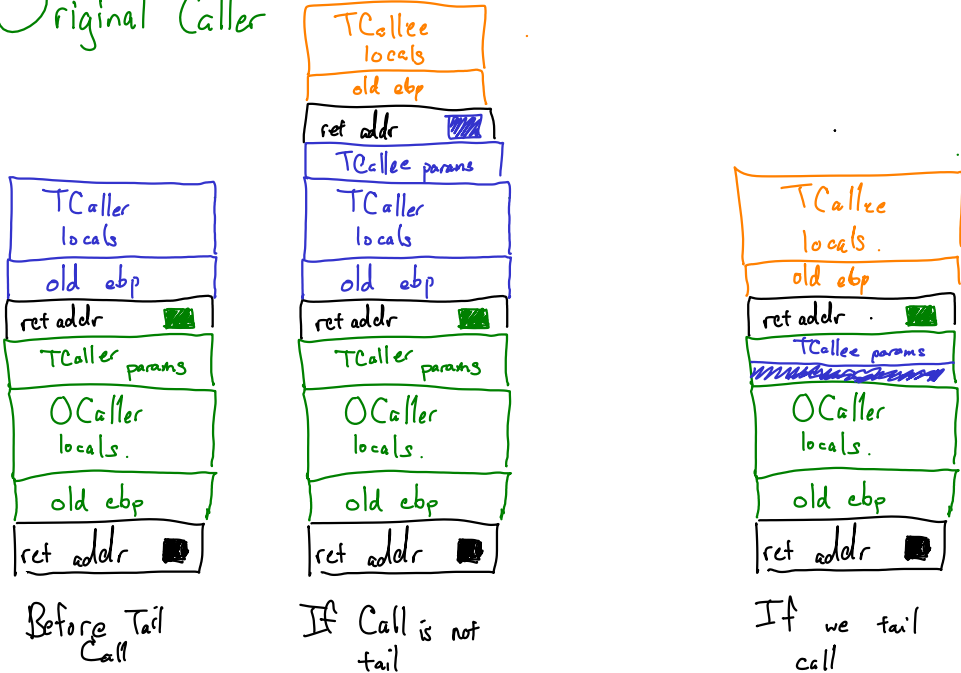


Tail Callee

Tail Caller

Original Caller

$$T_{\text{Caller}} \# \text{params} \geq T_{\text{Callee}} \# \text{params}$$

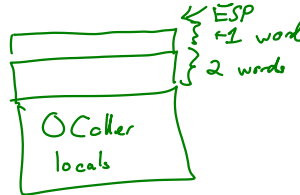


- Replace my own params w/ params of TCallee
- Tear down my stack (exit 1/2 of C callee) (pop "old ebp" into EBP)
- Imp (NOT call) the TCallee fn

Hypothetical: TCaller has 2 params & TCallee has 3



Just normal call



TCO info

- Is this a tail pos?  
Static info (compile-time)
- $T_{\text{Callee}} \# \text{params} \geq T_{\text{Caller}} \# \text{params}$   
Dynamic info (run-time).

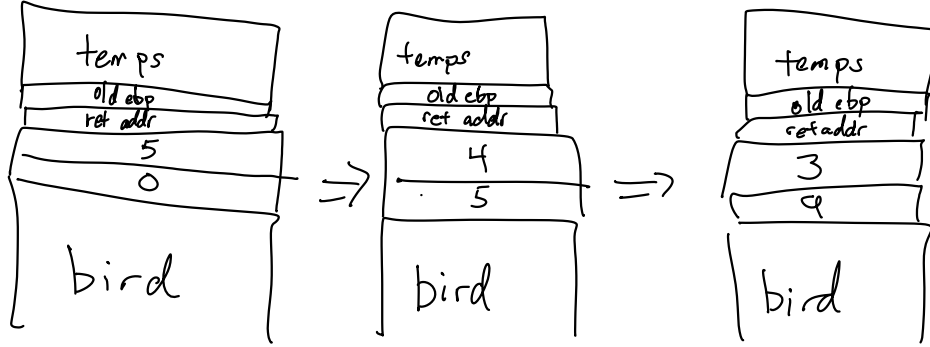
```

def f x y =
  x - y
end
def g x =
  x
end
def h x b q =
  (if b then f2 else q) x
end
  
```

```
def summate n acc =
  if n=0 then acc else
    summate (n-1) (acc+n)
end
```

summate 5 0

```
int summate (int n, int acc) {
  while (n>0) {
    acc+=n;
    n-=1;
  }
  return acc;
}
```



# TCallee params > # TCaller params

Callee:                    set up locals            tear down locals            tear down params

Caller: set up params

