

TF<sub>b</sub>R

$$\text{destruct record} \quad \text{construct record}$$

$$(\text{Function } a : \{q : \text{Int}\} \rightarrow a.q) \quad \{q = 5\} \Rightarrow 5$$

Evaluation Relations  $e \Rightarrow v$

Type Relation  $\vdash e : \tau$

Var  
Proj  
Fun  
App

$$\{a : \{q : \text{Int}\}\} \vdash a : \{q : \text{Int}\}$$

$$\{a : \{q : \text{Int}\}\} \vdash a.q : \text{Int}$$

$$\emptyset \vdash (\text{Function } a : \{q : \text{Int}\} \rightarrow a.q) : \{q : \text{Int}\} \rightarrow \text{Int}$$

$$\emptyset \vdash (\text{Function } a : \{q : \text{Int}\} \rightarrow a.q) \quad \{q = 5\} : \text{Int}$$

$$\text{Rec} \frac{\text{Int} \quad \emptyset \vdash 5 : \text{Int}}{\emptyset \vdash \{q = 5\} : \{q : \text{Int}\}}$$

Appl

$$\frac{\vdash e_a : \tau \rightarrow \tau' \quad \vdash e_a : \tau}{\vdash e_a e_a : \tau'}$$

Fun

$$\frac{\Gamma, x : \tau \vdash e : \tau'}{\vdash (\text{Function } x : \tau \rightarrow e) : \tau \rightarrow \tau'}$$

Proj

$$\frac{\vdash e : \{l_1 : \tau_1, \dots, l_n : \tau_n\}, l = l_k}{\vdash e.l : \tau_k}$$

$$\frac{\vdash e_1 : \tau_1, \dots, \vdash e_n : \tau_n}{\vdash \{l_1 = e_1, \dots, l_n = e_n\} : \{l_1 : \tau_1, \dots, l_n : \tau_n\}}$$

F<sub>b</sub>R

$$(\text{Function } a \rightarrow a.q) \quad \{q = 5\} \Rightarrow 5$$

$$(\text{Function } a \rightarrow a.q) \quad \{q = 5, z = \text{True}\} \Rightarrow 5$$

TF<sub>b</sub>R

w.s.t.  $\emptyset \vdash (\text{Function } a : \{q : \text{Int}\} \rightarrow a.q) \quad \{q = 5, z = \text{True}\} : \text{Int} \leftarrow \text{This claim does not hold}$

*Don't match*  $\uparrow$  has type  $\{q : \text{Int}, z : \text{Bool}\}$

This is disappointing b/c the record I have contains everything I need and more!

## Subtyping

"A type is a set of values"  
"A subtype is a subset of values"

Eval  $e \Rightarrow v$   
Typing  $\Gamma \vdash e : \tau$   
Subtyping  $\tau_1 \leq \tau_2$

Grammar:  $\tau \leq \tau$

$\leq$   $\leq$

Meaning:  $\tau_1 \leq \tau_2$  claims that all values in  $\tau_1$  are also in  $\tau_2$

Intuition: If  $\tau_1 \leq \tau_2$  then  $\tau_1$  can be used anywhere that  $\tau_2$  is expected.

TFbR



Reflexivity

$$\frac{}{\tau \leq \tau}$$

Transitivity

$$\frac{\tau_1 \leq \tau_2 \quad \tau_2 \leq \tau_3}{\tau_1 \leq \tau_3}$$

Symmetry

$$\frac{\tau_1 \leq \tau_2}{\tau_2 \leq \tau_1} \quad \text{...}$$

Record

$$\frac{\tau_1 \leq \tau'_1 \quad \dots \quad \tau_n \leq \tau'_n}{\{\ell_1 : \tau_1, \dots, \ell_n : \tau_n\} \leq \{\ell'_1 : \tau'_1, \dots, \ell'_n : \tau'_n\}}$$

Int < Int

$$\{q : \text{Int}, z : \text{Bool}\} \leq \{q : \text{Int}\}$$

$$\{q : \text{Int}, y : \text{Int}, z : \text{Bool}\} \leq \{q : \text{Int}, z : \text{Bool}\}$$

Write a

record rule

$$\{d : \{q : \text{Int}, z : \text{Bool}\}\} \leq \{d : \{q : \text{Int}\}\}$$

Function r:  $\{d : \{q : \text{Int}\}\} \rightarrow r.d.q$

More fields = "width subtyping"

More specific fields = "depth subtyping"

STFbR

$$\frac{\Gamma \vdash e : \tau \quad \tau \leq \tau'}{\Gamma \vdash e : \tau'}$$

$$\frac{\emptyset \vdash 5 : \text{Int} \quad \emptyset \vdash \text{False} : \text{Bool}}{\emptyset \vdash \{a=5, b=\text{False}\} : \{a : \text{Int}, b : \text{Bool}\}}$$

$$\frac{\emptyset \vdash 5 : \text{Int} \quad \emptyset \vdash \text{False} : \text{Bool}}{\emptyset \vdash \{a=5, b=\text{False}\} : \{a : \text{Int}\}}$$

# Subtyping Functions

STFbR

$\tau ::= \text{Int} \mid \text{Bool} \mid \{\ell : \tau, \dots\} \mid \tau \rightarrow \tau$

(Function  $f : \{a : \text{Int}\} \rightarrow \{b : \text{Int}\} \rightarrow f \{a = b\}$ )

(Function  $r : \{a : \text{Int}\} \rightarrow \{b = r.a, c = \text{False}\}$ )  
 $\{a : \text{Int}\} \rightarrow \{b : \text{Int}, c : \text{Bool}\}$

Focusing on output:

(Not the whole picture)

$$\frac{\tau_1 \leq \tau_2}{\tau \rightarrow \tau_1 \leq \tau \rightarrow \tau_2}$$

With input (the whole picture)

$$\text{Func Sub } \frac{\tau'_2 \leq \tau'_1 \quad \tau_1 \leq \tau_2}{\tau'_1 \rightarrow \tau'_1 \leq \tau'_2 \rightarrow \tau_2}$$

"contravariance on input"