

Logic

\forall

for all

\exists

there exists.

I+ is cloudy..

I am flying.

$\forall \text{day} \in \text{days. } \text{cloudy}(\text{day})$

$\frac{\text{foo (bar)}}{\text{predicate}}$

$P(n)$

Quantifier Ordering

$\exists x. \forall y. x > y$	there is a number which is bigger than every number	FALSE
$\forall y. \exists x. x > y$	for each number, some number is bigger than it	TRUE

$\forall n \in S. P(n)$

$\forall n. \underset{\text{and}}{\wedge} n \in S P(n)$

and

$\underset{\text{or}}{\vee}$

or

If I have no work and it is the weekend then I will sleep in.

if $\text{work} = \emptyset \wedge \text{weekend}(\text{today})$ then sleepIn(me)

$\text{work} = \emptyset \wedge \text{weekend}(\text{today}) \Rightarrow \text{sleepIn(me)}$

not going to use this for implication*

$$\frac{\text{work} = \emptyset \quad \text{weekend}(\text{today})}{\text{sleepIn(me)}}$$

Strings composed of $\Delta \square \wedge$

Inference Rules:

$$1 \frac{\wedge}{\wedge \square} \quad 2 \frac{\square}{\square \square} \quad 3 \frac{\wedge \square}{\wedge \Delta} \quad 4 \frac{\wedge \square}{\wedge \Delta \square} \quad 5 \frac{\wedge \Delta}{\Delta \Delta} \quad 6 \frac{}{\wedge}$$

↓ axiom (no premises)

Prove $\Delta \Delta$

\wedge therefore $\wedge \square$; therefore $\wedge \Delta$ therefore $\Delta \Delta$

$$\frac{6 \frac{\wedge}{\wedge}}{1 \frac{\wedge \square}{\wedge \square}} \frac{3 \frac{\wedge \square}{\wedge \Delta}}{5 \frac{\wedge \Delta}{\Delta \Delta}}$$

Metavariable

$s ::= \text{strings containing } \wedge, \Delta, \square \text{ and nothing else}$

$$1 \frac{\wedge}{\wedge \square \square} \quad 2 \frac{s \square}{s \Delta} \quad 3 \frac{s \square}{s \Delta \Delta} \quad 4 \frac{\wedge s}{\square s} \quad 5 \frac{\square s \Delta}{s} \quad 6 \frac{}{\wedge}$$

Prove " \wedge ".

Prove " \square ".

$$\begin{array}{c}
 6 \frac{\wedge}{\wedge} \\
 1 \frac{\wedge}{\wedge \square \square} \\
 3 \frac{\wedge \square \square}{\wedge \square \Delta \Delta} \\
 4 \frac{\wedge \square \Delta \Delta}{\square \square \Delta \Delta} \\
 s = \epsilon \quad 5 \frac{\square \square \Delta \Delta}{\epsilon} \\
 \hline
 s = \epsilon \quad 6 \frac{\wedge}{\square}
 \end{array}$$

$s = " \wedge \square "$
 $s = " \square \Delta \Delta "$
 $s = " \square \square \Delta \Delta "$
 $s = \epsilon$

$$1 \frac{s}{s\square\square}$$

$$2 \frac{s_1 \not\sim s_2}{s_2 \Delta s_1 s_2}$$

$$3 \frac{s_1 \square \Delta s_2}{s_1 s_2}$$

$$4 \frac{s_2}{s_1 \not\sim s_2} \frac{s_2}{s_2}$$

$$5 \quad \not\sim$$

Prove: $\Delta \Delta \not\sim \square \square$

$$\begin{array}{c} 5 \overline{\not\sim} \\ 4 \overline{\not\sim} \\ 1 \overline{\not\sim} \\ 2 \overline{\not\sim} \\ 2 \overline{\not\sim} \\ \hline \Delta \Delta \not\sim \square \square \end{array}$$

$$\begin{array}{c} 5 \overline{\not\sim} \\ 2 \overline{\Delta \not\sim} \\ 2 \overline{\Delta \Delta \not\sim} \\ 4 \overline{\Delta \Delta \not\sim \square \square} \\ \hline \Delta \Delta \not\sim \square \square \end{array}$$

Relations

$$R \subseteq S \times T \quad \text{on } \mathbb{Z}^{\alpha}, \mathbb{Z}^{\alpha}, \dots$$

Relation $\subseteq \mathbb{Z}^{\alpha} \times \mathbb{Z}^{\alpha}$

" \leq "

$$\leq^{\alpha} = \left\{ (2, 4), (2, 5), \dots \right\}$$

$$\frac{n \in \mathbb{Z}^{\alpha}}{n \leq n}$$

$$\frac{n \leq m}{n \leq m+1}$$

$$\frac{n \leq m}{n-1 \leq m}$$