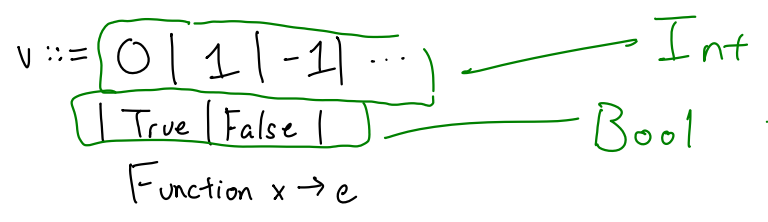


Types

A type is a set of values.



$\tau ::= \text{Int} \mid \text{Bool} \mid \dots$

$\Gamma \vdash e : \tau$

$e \Rightarrow v$

Assuming Γ , e has type τ .

Evaluating e should produce v

If e is evaluated, then either $e \Rightarrow v$ for $v \in \tau$ or e loops forever.

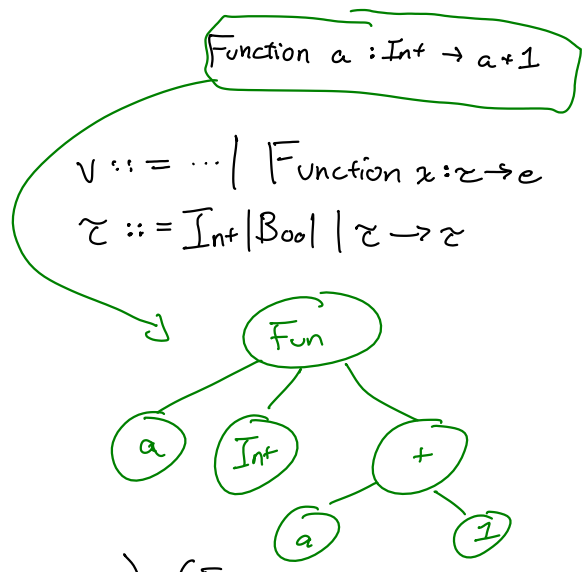
$\frac{\Gamma \vdash e_1 : \text{Bool} \quad \Gamma \vdash e_2 : \text{Bool}}{\Gamma \vdash e_1 \text{ And } e_2 : \text{Bool}}$

Function $n : \text{Int} \rightarrow_{n=0}$

$\frac{\Gamma, x : \tau \vdash e : \tau'}{\Gamma \vdash (\text{Function } x : \tau \rightarrow e) : \tau \rightarrow \tau'}$

$\frac{\Gamma \vdash e_1 : \tau_1 \rightarrow \tau_2 \quad \Gamma \vdash e_2 : \tau_1}{\Gamma \vdash e_1 e_2 : \tau_2}$

$\frac{x : \tau \in \Gamma}{\Gamma \vdash x : \tau}$



TFB
↑
no Let Rec
↓
cannot type Y-combinator

TcFb
↑
yes Let Rec

$(\text{Function } x : \tau \rightarrow x x) \quad (\text{Function } x : \tau \rightarrow x x)$
 $\tau_\omega = \tau_\omega \rightarrow \tau_a = \tau_\omega \rightarrow \tau_\omega \rightarrow \tau_2 = \dots$

TFB is normalizing

TFBR

$e ::= \dots \mid \{ l = e, \dots \} \mid e.l$
 $v ::= \dots \mid \{ l = v, \dots \}$
 $\tau ::= \dots \mid \{ l : \tau, \dots \}$

$\frac{\Gamma \vdash e_1 : \tau_1 \quad \dots \quad \Gamma \vdash e_n : \tau_n}{\Gamma \vdash \{ l_1 = e_1; \dots; l_n = e_n \} : \{ l_1 : \tau_1, \dots, l_n : \tau_n \}}$
 $\frac{\Gamma \vdash e : \{ l_1 : \tau_1, \dots, l_n : \tau_n \} \quad i \in \{ 1, \dots, n \}}{\Gamma \vdash e.l_i : \tau_i}$

