

• 1

(Function  $a \rightarrow a+a+a+a+a$ ) 1

(Function  $f \rightarrow$  Function  $g \rightarrow$  Function  $n \rightarrow f(g\ n)$ ) • (Function  $a \rightarrow a$ ) 5

$\rightarrow$  • ((Function  $a \rightarrow a$ ) 5)

$\xrightarrow{\sim}$  • 5

Soundness is an ex. of relating three trees:  $\Gamma \vdash e :: \tau$ ,  $e \Rightarrow v$ ,  $\Gamma \vdash v :: \tau$

$e ::= v \mid e \mid e \text{ Or } e \mid \text{Not } e$

$v ::= \text{True} \mid \text{False}$

Let  $e \xrightarrow{Fb} v$  denote the operational semantics of Fb. Let  $e \xrightarrow{Bool} v$  denote the operational semantics of BOOL. BOOL is a subset of Fb; that is, if  $e \xrightarrow{Bool} v$  then  $e \xrightarrow{Fb} v$ .  
 (Note: "given" under  $e \xrightarrow{Bool} v$ , "Want to show" under  $e \xrightarrow{Fb} v$ )

Proof

By induction on the height of the proof of  $e \xrightarrow{Bool} v$ .

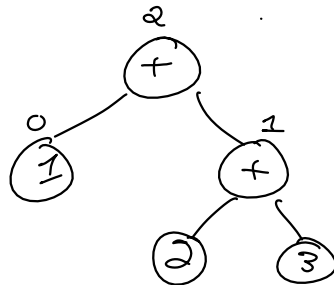
Proceed by case analysis on the rule used in  $e \xrightarrow{Bool} v$ .

Let  $P(k) \equiv$  If  $e \xrightarrow{Bool} v$  then  $e \xrightarrow{Fb} v$  for proof of  $e \xrightarrow{Bool} v$  with height  $k$ .

Case of Value Rule: If the Value Rule is used to prove  $e \xrightarrow{Bool} v$ , then  $e = v$ .

So by the Value Rule of Fb,  $v \xrightarrow{Fb} v$  and this case is finished.

$P(0), \dots, P(k)$  implies  $P(k+1)$



Normalizing: every expression evaluates to at least one value  
 $\forall e. \exists v. e \Rightarrow v$

# Recursion

Let countdown' = Function self  $\rightarrow$  Function n  $\rightarrow$   
If n=0 Then 0 Else self self (n-1)

In

Let countdown = countdown' countdown' In  
countdown 5

$\lambda$  (Function recurse  $\rightarrow$  Function n  $\rightarrow$  ... recurse (n-1))

$\lambda 2$  (Fun f1  $\rightarrow$  Fun f2  $\rightarrow$  Fun n  $\rightarrow$  ...)

(Fun f1  $\rightarrow$  Fun f2  $\rightarrow$  Fun n  $\rightarrow$  ...)

Let countdown' = Function self  $\rightarrow$  Function n  $\rightarrow$   
If n=0 Then 0 Else self self (n-1) C

In

Let countdown = countdown' countdown' In  
countdown 5

$\downarrow$

Let countdown = C C In  
countdown 5

$\downarrow$

(C C) 5

(Function n  $\rightarrow$  If n=0 Then 0 Else C C (n-1)) 5

$\downarrow$

If 5=0 Then 0 Else C C (5-1)

$\downarrow$

(C C) (5-1)

Let somefun = Function recurse  $\rightarrow$  Function x  $\rightarrow$  ... In

Let Ycomb = Function f  $\rightarrow$

Let recuser = Function self  $\rightarrow$  Function arg  $\rightarrow$   
f (self self) arg

In

f (recuser recuser)

In

Ycomb somefun 5