

Ways to define behavior:

Write rules
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Encoding

This week

Records

\boxed{FbR}

$e ::= \underbrace{\dots}_{Fb} \mid \underbrace{\{l = e, \dots\}}_{\substack{\text{? zero or more?} \\ \text{one or more?}}} \mid e.l$

$v ::= \underbrace{\dots}_{Fb} \mid \{l = v, \dots\}$

$l ::= (\text{identifiers})$

In all record expressions $\{l_1 = e_1, \dots, l_n = e_n\}$,
all labels l_k must be unique.

Record expression elements are unordered:

Record

Projection

$$\begin{aligned} \{a=4, b=\text{True}\} &\Rightarrow \{a=4, b=\text{True}\} \\ \{a=1+2, b=\text{False}\}.a &\Rightarrow 3 \\ (\{a=\text{Function } x \rightarrow x\}.a) \ 4 &\Rightarrow 4 \\ 5.a &\not\Rightarrow \\ \cancel{\{a=4, a=3\}.a} &\Rightarrow 4 \\ \cancel{\{a=4, a=3\}.a} &\Rightarrow 3 \\ \forall i \in \{1, \dots, n\}. e_i &\Rightarrow v_i \\ \frac{e_1 \Rightarrow v_1 \quad \dots \quad e_n \Rightarrow v_n}{\{l_1 = e_1, \dots, l_n = e_n\} \Rightarrow \{l_1 = v_1, \dots, l_n = v_n\}} \\ \frac{e \Rightarrow \{l_1 = v_1, \dots, l_n = v_n\} \quad l = l_i}{e.l \Rightarrow v_i} \end{aligned}$$

Variants

Fb Exprs

$e ::= \dots | \lambda e | \text{Match } e \text{ With } \{\lambda x \rightarrow e\} \dots$

$v ::= \dots | \lambda v$

Fb vals

Variant

$$e \Rightarrow v$$

$$\lambda e \Rightarrow \lambda v.$$

$$\text{'Foo } 5 \Rightarrow \text{'Foo } 5$$

$$\text{'Bar (True And False)} \Rightarrow \text{'Bar False}$$

Match 'Bar 3 With

$$\{\text{'Foo } a \rightarrow a+1\}$$

$$\{\text{'Bar } b \rightarrow b-1\}$$

$\Rightarrow 2$

Match

$$e_0 \Rightarrow \lambda_i v_i$$

$$e_0 \Rightarrow \lambda_i v_i \quad \lambda = \lambda_i \quad e_i[v/x_i] \Rightarrow v'$$

$$\text{Match } e_0 \text{ With } \{\lambda_1 x_1 \rightarrow e_1 | \dots | \lambda_n x_n \rightarrow e_n\} \Rightarrow v'$$

Records & Variants: Dual Encoding

Can encode FbV in FbR because functions!
 Can encode FbR in FbV

Given a program in FbV , assign all labels ℓ a number. (e.g. $\text{Foo}=1, \text{Bar}=13, \text{etc.}$)

$$\ell e \stackrel{\text{def}}{=} \{ \text{type} = \#, \text{value} = e \}$$

Match e_0 With
 $\begin{cases} \ell_1 x_1 \rightarrow e_1 \\ \vdots \\ \ell_n x_n \rightarrow e_n \end{cases}$ $\stackrel{\text{let}}{=}$ Let $r = e_0$ In
 If $r.\text{type} = \#_1$ Then
 Let $x_1 = r.\text{value}$ In e_1
 Else If ...
 Else If ...
 Let $x_n = r.\text{value}$ In e_n
 Else
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Given a program written in FbR , assign all labels ℓ a number.

$$\{ \ell_1 = e_1, \dots, \ell_n = e_n \} \stackrel{\text{def}}{=} \left(\begin{array}{l} \text{Let } a_1 = e_1 \text{ In} \\ \quad \text{Let } a_n = e_n \text{ In} \\ \quad \text{Function } \text{lbl} \rightarrow \\ \quad \quad \text{If } \#_1 \text{ Then } a_1 \text{ Else} \\ \quad \quad \quad \vdots \\ \quad \quad \text{If } \#_n \text{ Then } a_n \text{ Else} \\ \quad \quad \quad \text{○ ○} \\ \end{array} \right)$$

$e.l \stackrel{\text{def}}{=} e \#$