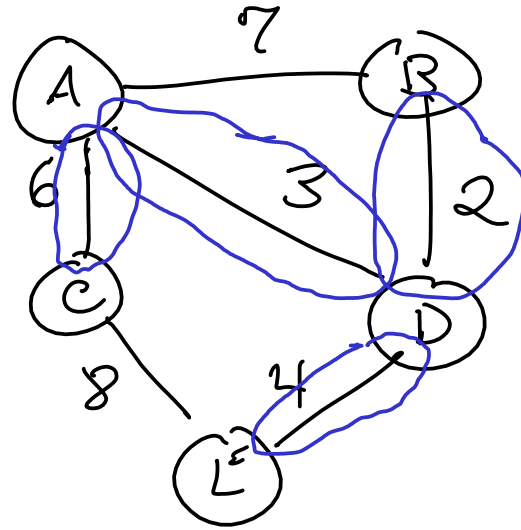
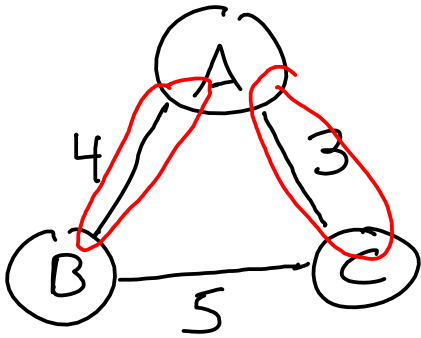
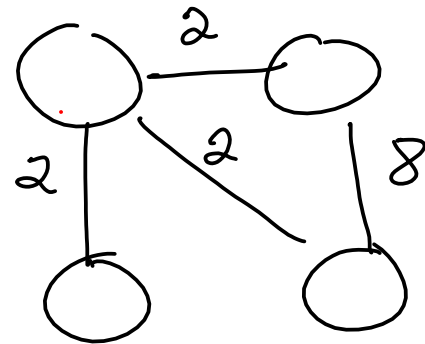
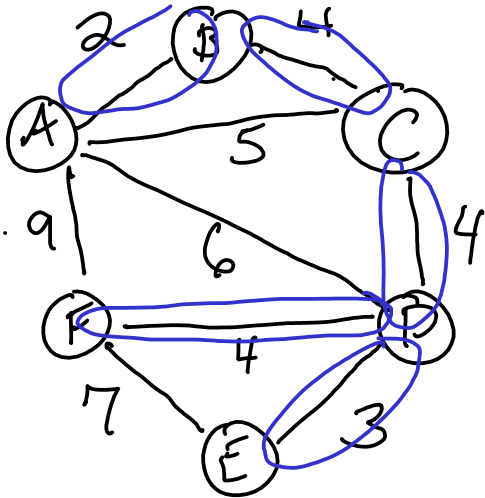


Minimum Spanning Trees

MST : set of edges which, given V in G , connect that graph, which $\subseteq E$, least cost

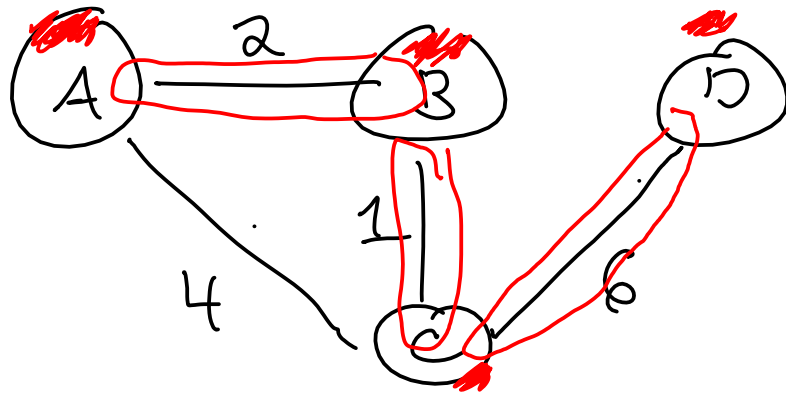


Start w/ cheapest edge



Prim's Algorithm

1. Start w/ cheapest edge in graph, add that
2. Remember: been to that edge's source & target
3. Find from edges leaving vertices to which you have been the cheapest that goes somewhere new
4. Add that edge; remember new place
5. Repeat 3 & 4 until connected. (been everywhere)



Prim's (more detail)

Function Prim's(g):
visited \leftarrow new Set
edges \leftarrow new Set
 $e \leftarrow$ cheapest edge in g
edges.add(e)

visited.add(e .source)
visited.add(e .target)

$q \leftarrow$ new PriorityQueue

for each edge leaving e .source or e .target:

if edge \notin edges:
q.enqueue(edge.weight, edge)

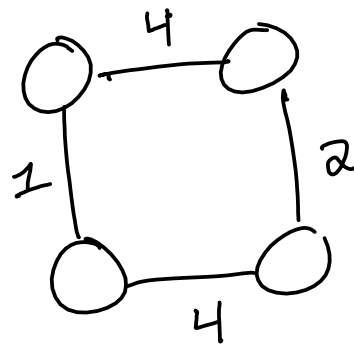
while visited $\neq g.V$

\rightarrow while q not empty:

$e \leftarrow$ q.dequeue()

if e .target \notin visited or e .source \notin visited:
edges.add(e)

if e .source \notin visited:
for each edge leaving e .source
 \rightarrow q.enqueue(edge.weight, edge)
if e .target \notin visited:
for each edge leaving ...
visited.add(e .source)
visited.add(e .target)



Kruskal's Algorithm

$g \leftarrow$ new Prio Queue
for every edge in g :
 $g.\text{enqueue}(\text{edge.weight}, \text{edge})$
 $\text{answer} \leftarrow$ new Set
while g not empty :
 $e \leftarrow g.\text{dequeue}$
 if e connects two prev disconnected subgraphs :
 $\text{answer.add}(e)$

Union-Find

