

ADT — abstract data type — List

data structure — Linked List

Dictionary ADT

2 type parameters K (key - lookup)
 V (value - result)

like "T" from lists

void insert(K key, V value) ← new mappings

void update(K key, V value) ← changing mappings

void remove(K key)

List $\langle K \rangle$ getKeys()

V get(K key)

Guarantee:

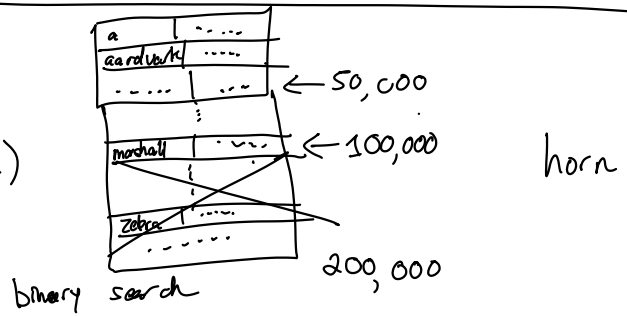
for every mapping, K is unique

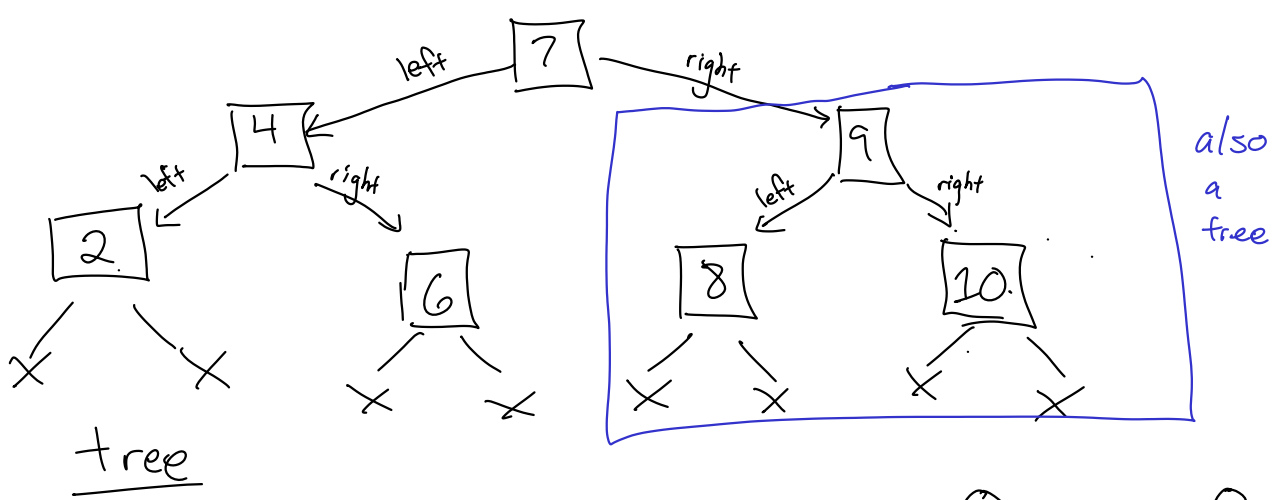
Array Dictionary

uses an array of pairs of K and V

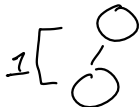
get(K) — binary search $O(\log n)$

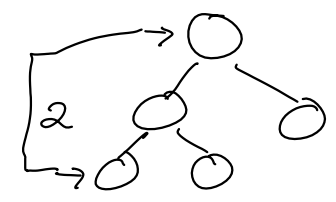
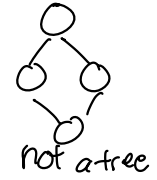
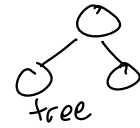
insert(K, V) — $O(n)$





Vocab

- Node — container holding data; point in the tree
- Child — a node to which this node points
- Parent — the node pointing to this one (of which there is at most one)
- Root — the node in a tree with no parent
- Leaf — a node w/ no children
- Ancestor / Descendant — $-1 \leq \emptyset \leq 1$ 
- Size — # nodes in tree
- Height — # of links between root & a deepest leaves



Tree: collection of nodes s.t.

- there is a root which has no parent
- all other nodes have a unique parent
- all nodes have the root as an ancestor.

or

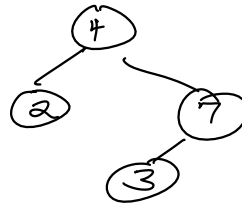
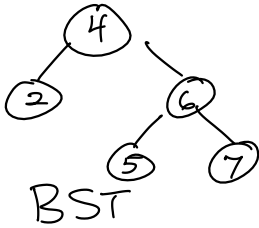
the empty tree

Binary Trees:

Trees in which each node has at most 1 left child and at most 1 right child

Binary Search Tree:

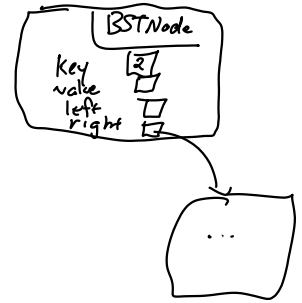
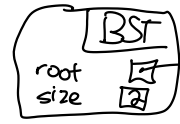
Every node to the left of a node N has key $<$ the key of N and every node to the right of a node N has key $>$ the key of N



Binary tree
not
BST

BST data structure implements dictionary ADT interface

Method `getSize()` :
Return `sizeOfTree(this → root)`
End Method



Method `sizeOfTree(Node* n)` :
If `n == nullptr` :
Return 0
Else :
 $sL \leftarrow \text{sizeOfTree}(n \rightarrow \text{left})$
 $sR \leftarrow \text{sizeOfTree}(n \rightarrow \text{right})$
 Return $sL + sR + 1$
End If
End Method

