

Merge Sort

- Split data into two piles
- Sort (recursively)
- Merge

Bubble sort

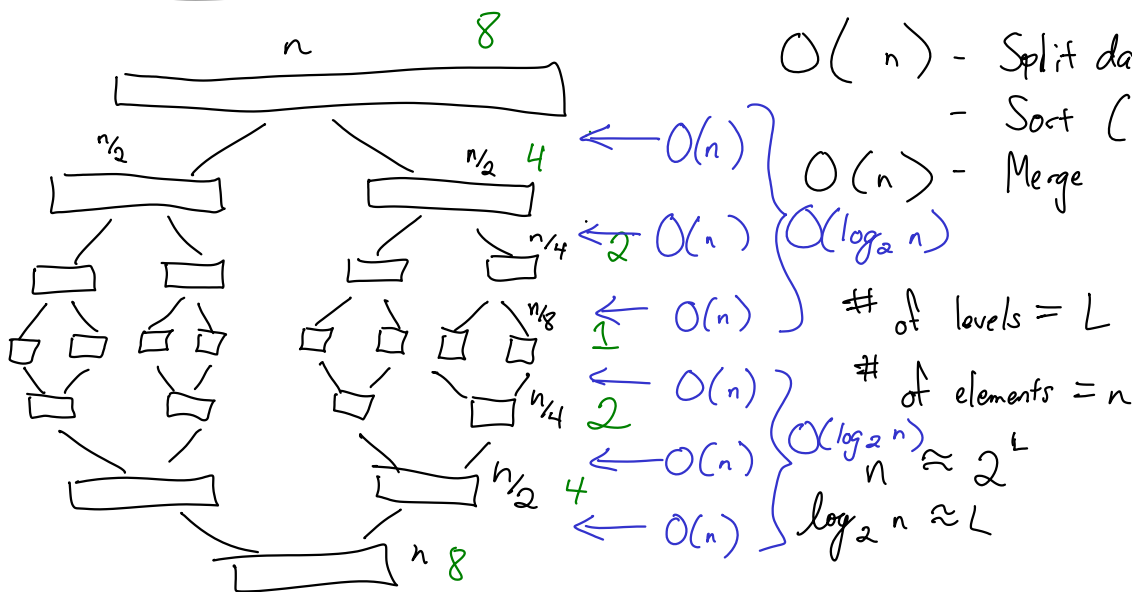
Insertion sort

Selection sort

$$O(n^2)$$

Recursive function

- Base case
- Inductive/recursive case (smaller problem)



$O(n)$ - Split data into two piles

- Sort (recursively)

$O(n)$ - Merge

$$O(\log_2 n)$$

of levels = L

of elements = n

$$n \approx 2^L$$

$$\log_2 n \approx L$$

$$O(n \log n)$$

Suppose $n = 1000000$. $n^2 = 1000000000000$
 $n \log_2 n = 20000000$

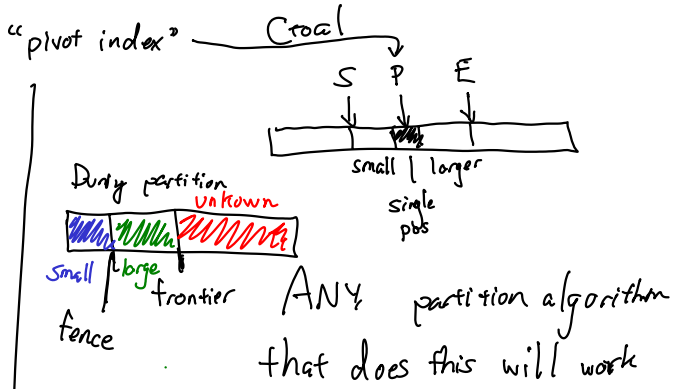
Quick Sort

- Recursive sort
 - Split array into two segments
 - s.t. all smaller things in 1st segment
 - all bigger things in 2nd segment
 - recursively sort segments
- In-place (no extra memory*)

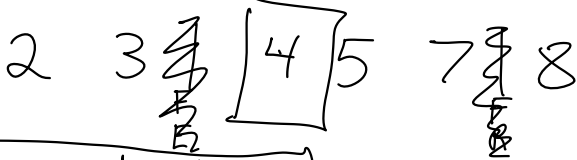
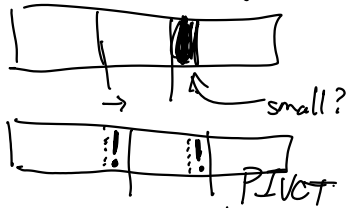
Lomuto
 Partition (A, start, end) ← returns "pivot index" inclusive

```

pivotIndex ← end
pivot ← A[pivotIndex]
fence ← start
frontier ← start
While frontier < end :
    If A[frontier] ≥ pivot :
        ⌊ (large element, so we're okay)
    Else :
        swap A[fence] w/ A[frontier]
        fence ← fence + 1
    End If
    frontier ← frontier + 1
End While
swap A[pivotIndex] w/ A[fence]
Return fence
End Function
    
```



At each step, one more element is in the right place (on the correct side of pivot) starting at left and going to right.



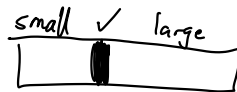
Hoare algorithm

Not discussing this one

Function QS(A, start, end):

```

size ← end - start + 1
If size < 2 :
    Return
pivotIndex ← partition (A, start, end)
QS (A, start, pivotIndex - 1)
QS (A, pivotIndex + 1, end)
End Function
    
```



3 4 | 9 8 9 7 | 5

F
E

F
R

pivot

QS complexity: worst case $O(n^2)$

Randomized QS

- At start of partition,
pick a random index
& swap w/ last index

	QS	RQS
worst case	$O(n^2)$	$O(n^2)$
expected worst case	$O(n^2)$	$O(n \log n)$

	bad user	bad luck
worst	✓	✓
expected worst	✓	~