

Pseudocode

1000 sets + 2000 gets + 1000 compares
+ 1 return

Function isSorted (array, size) :

For i From 0 to size - 2 :

If array[i] > array[i+1] : } 1 step

Return False

Return True

End Function

} O(size)

Function isSorted (array, size)

For i From 0 to size - 1 :

For j From i+1 to size - 1 :

If array[i] > array[j] : } 1 step

Return False

Return True

End Function

size = n

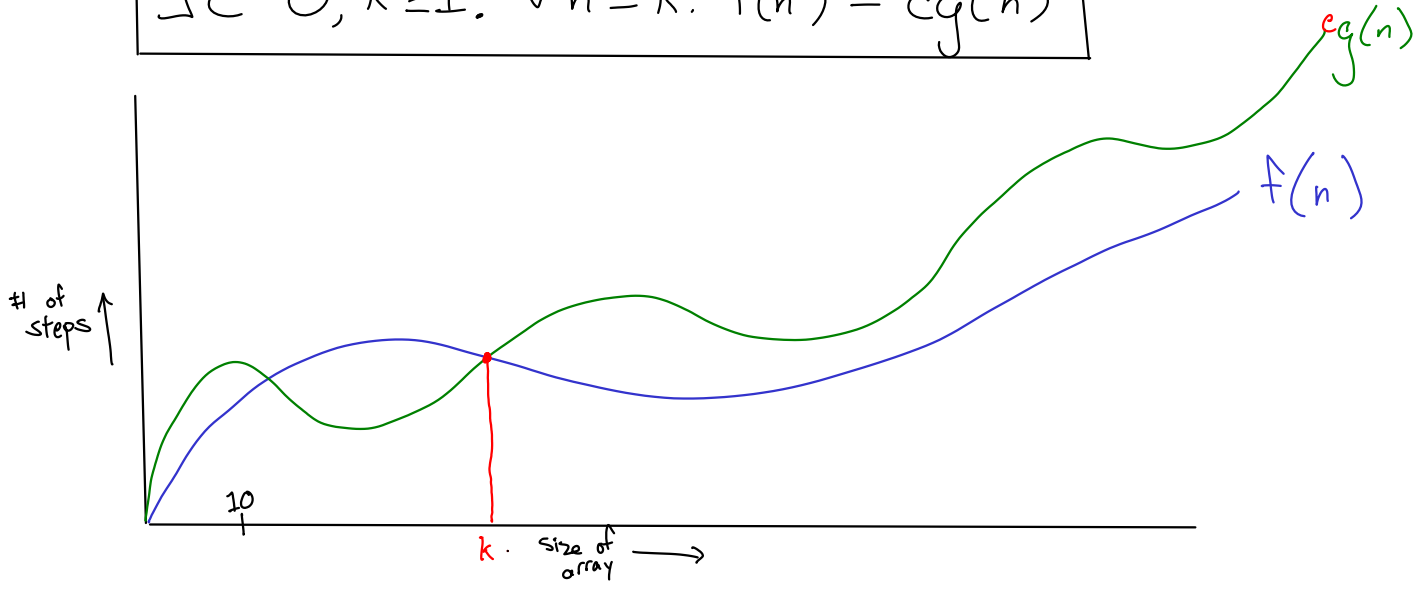
$$\left. \begin{array}{l} 0+1+2+\dots+n \\ \sum_{i=1}^n i = \frac{n(n+1)}{2} \end{array} \right\}$$

$$\boxed{\frac{n^2+n}{2} \text{ is } O(n^2)}$$

$O(n)$ = "proportional to n "

$f(n)$ is $O(g(n))$

$$\exists c > 0, k \geq 1. \forall n \geq k. f(n) \leq cg(n)$$



$$\frac{n^2+n}{2} \text{ is } O(n^2)$$

$f(n)$

$g(n)$

$$\exists c > 0, k \geq 1. \forall n \geq k. \frac{n^2+n}{2} \leq cn^2$$

Let $c=5$. Then it is enough to show that

$$\exists k \geq 1. \forall n \geq k. \frac{n^2+n}{2} \leq 5n^2$$

Let $k=10$. Then it is enough to show that

$$\forall n \geq 10. \frac{n^2+n}{2} \leq 5n^2$$

$$\forall n \geq 10. \frac{n^2}{2} + \frac{n}{2} \leq 5n^2$$

$$\forall n \geq 10. \frac{n}{2} \leq \frac{9}{2}n^2$$

$$\forall n \geq 10. n \leq 9n^2$$

$$\forall n \geq 10. 1 \leq 9n$$

$$\frac{n^2+n}{2} \text{ is } O(n^2)$$

$$\exists c > 0, k \geq 1. \forall n \geq k. \frac{n^2+n}{2} \leq cn^2$$

$$\forall n \geq 2. 2 \leq n \quad \text{a positive}$$

$$\forall n \geq 2. 2n \leq n^2 \quad \leftarrow$$

$$\forall n \geq 2. \frac{n}{2} \leq \frac{n^2}{4}$$

$$\forall n \geq 2. 1 \leq 1$$

$$\forall n \geq 2. n^2 \leq n^2$$

$$\forall n \geq 2. \frac{n^2}{2} \leq \frac{n^2}{2}$$

If $a \leq c$
and $b \leq d$
then $a+b \leq c+d$

$$\forall n \geq 2. \frac{n}{2} + \frac{n^2}{2} \leq \frac{n^2}{4} + \frac{n^2}{2}$$

$$\forall n \geq 2. \frac{n^2+n}{2} \leq \frac{3n^2}{4}$$

$$\exists c > 0. \forall n \geq 2. \frac{n^2+n}{2} \leq cn^2$$

$$\exists c > 0, k \geq 1. \forall n \geq k. \frac{n^2+n}{2} \leq cn^2$$

$$\frac{n^2+n}{2} \text{ is } O(n^2)$$