

# Graph Algorithms

Graph: Set of vertices  
Set of edges

edge: source vertex  
target vertex  
weight  
label

alg	in	out	ds
DFS	Graph start $V$ end $V$	bool	Stack
BFS	Graph start $V$ end $V$	path	Queue
BFS All	Graph start $V$	Dictionary end $V \mapsto \#$ steps	Queue

Function BFSAll(Graph  $g$ ,  $V_{start}$ ):

Dictionary  $\langle V, \text{int} \rangle$   $dist \leftarrow$  new HT

$dist.insert(start, 0)$

Queue  $\langle V \rangle$   $q \leftarrow$  new LQ

$q.enqueue(start)$

While  $q$  is not empty:

$V_{current} \leftarrow q.dequeue()$

$int\ curCost \leftarrow dist.get(current)$

For each neighbor  $n$  of  $current$ :

If  $n$  is not a key in  $dist$ :

$q.enqueue(n)$

$dist.insert(n, curCost + 1)$

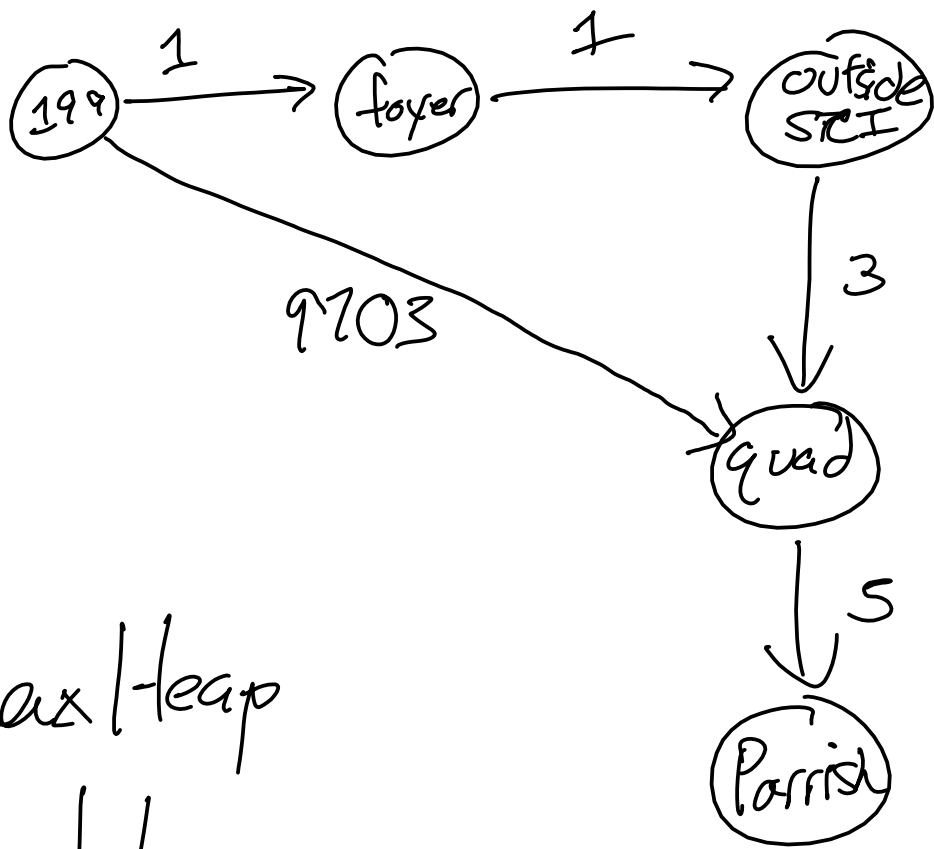
Return  $dist$

$O(|V|)$

$O(|V|)$

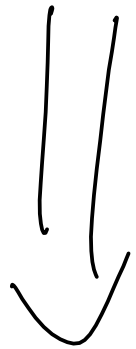
$O(1)$

$O(|E|)$



outside SCI -2  
 quad -9703

MaxHeap



MinHeap

negate  
all

priorities —

Function Dijkstras(Graph g, V start) :

Dictionary  $\langle V, W \rangle$  cost  $\leftarrow$  new HT  
cost.insert(start, 0)

PQ  $\langle W, V \rangle$  pq  $\leftarrow$  new MinHeap

pq.enqueue(0, start)

While pq is not empty :

V current  $\leftarrow$  pq.dequeue()

W curCost  $\leftarrow$  cost.get(current)

For each outgoing edge e of current :

V n  $\leftarrow$  e.target

W newCost  $\leftarrow$  curCost + e.weight

IF n is not a key in cost :

*If not in cost*  
cost.insert(n, newCost)

pq.enqueue(newCost, n)  $\leftarrow$  min heap

Else If

newCost  $<$

cost.get(n) :

cost.update(n, newCost)

pq.~~enqueue~~(newCost, n)  
*update*

$O(|V|^2 \log(|V|))$

Other versions

$O(|V|^2)$

$O(|V|)$

$O(|E|)$

Return cost

$O(|V|)$