

HashTable is a Dictionary

all operations to  $O(1)$  (on average, and w/ good hash)

1. Storing k/v pairs in an array
2. Hashing fn:  $k \rightarrow \text{int}$
3. mod int to be in array bounds
4. Deal w/ "collisions"

# Terms

size — # of mappings

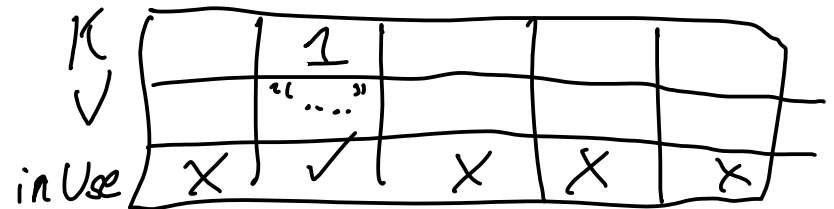
capacity — total allocated size of array

# Collision Resolution

Linear probing (Tuesday)

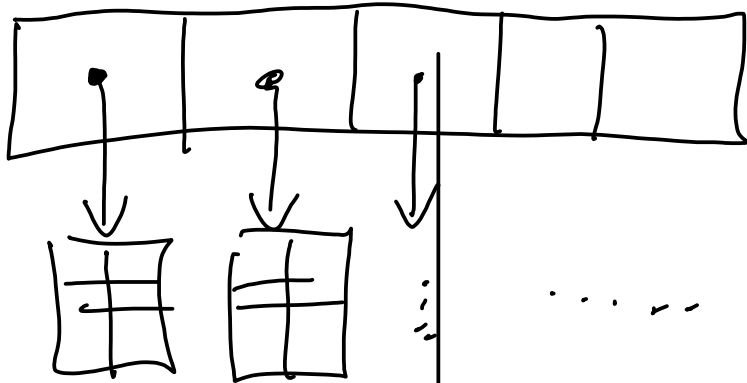
if in use, go next slot

0 1 2  
↓ ↓ ↓

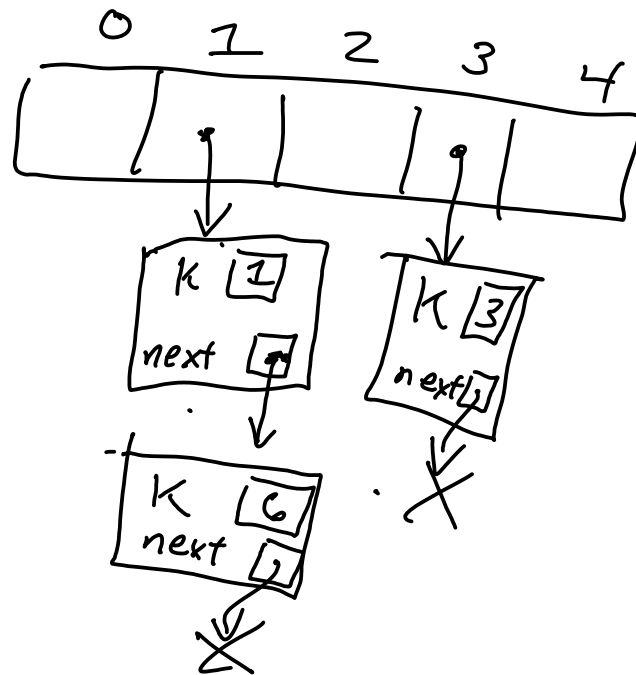


Other strategy:

# Chaining



insert (1, ...)  
insert (6, ...)  
insert (3, ...)  
insert (11, ...)



Dictionary impl:  
Linear Dictionary:  
list (unsorted) of  
K/V pairs

# Expanding HT Capacity

load factor:  $\frac{\text{size}}{\text{Capacity}}$

max load factor

small  
MLF  
< collisions  
> mem usage



large  
MLF  
> collisions  
< memory usage

---

Method `expandCapacity()`:

`newArray` ← new array of K/V pair lists; 2x size of prev

for each k/v pair in `oldArray`:

`h` ← hash key

`i` ← `h % newArray capacity`

`newArray[i].insertAtTail(pair)`

update capacity

update array (to be `newArray`)

hash functions!

$K \rightarrow \text{int}$

```
int hash(string s) {  
    return 0;  
}
```

good hash:  
distribute  $K$  evenly  
over int

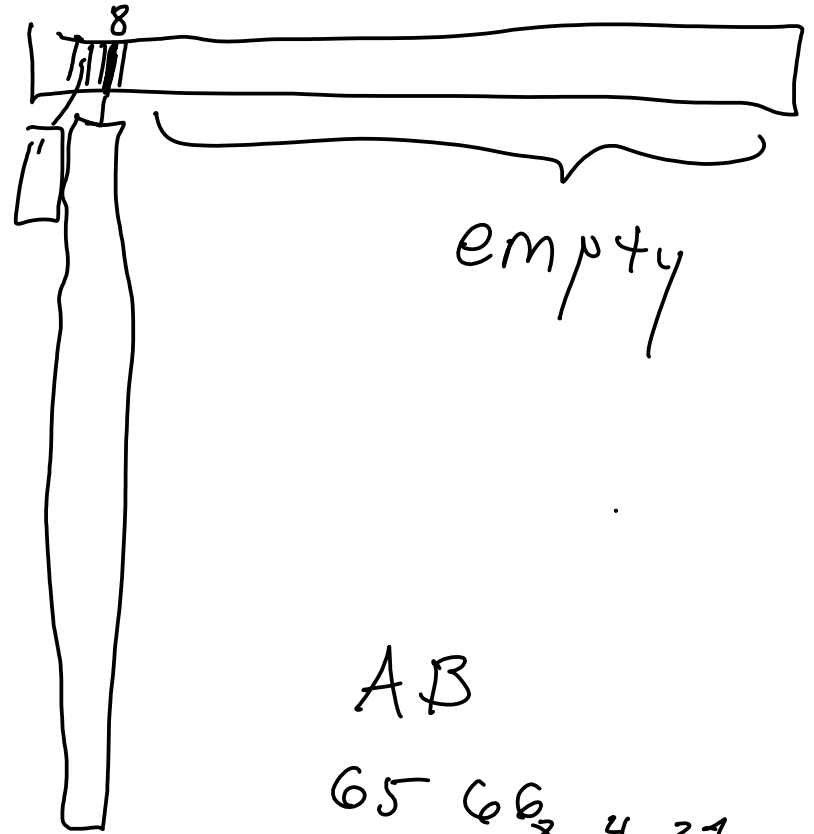
↑  
NOT A GOOD HASH

```

int hash(string s) {
    return s.length();
}

```

Better, but that's not saying much



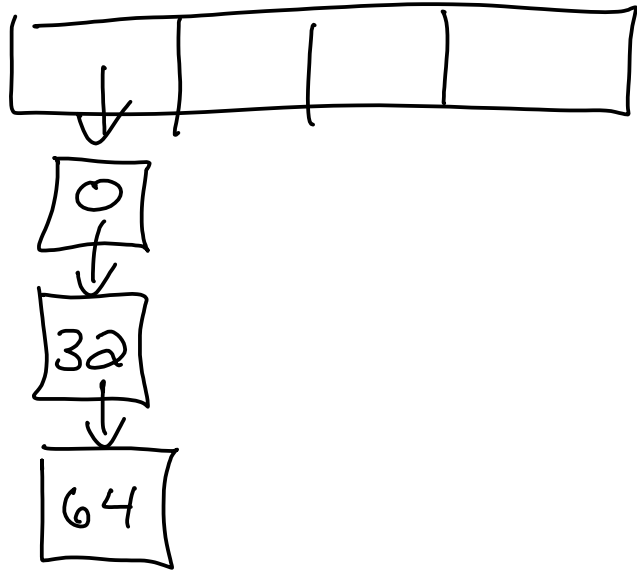
```

int hash(string s) {
    int n = 0;
    for (int i = 0; i < s.length(); i++) {
        n *= 31;
        n += s[i];
    }
    return n;
}

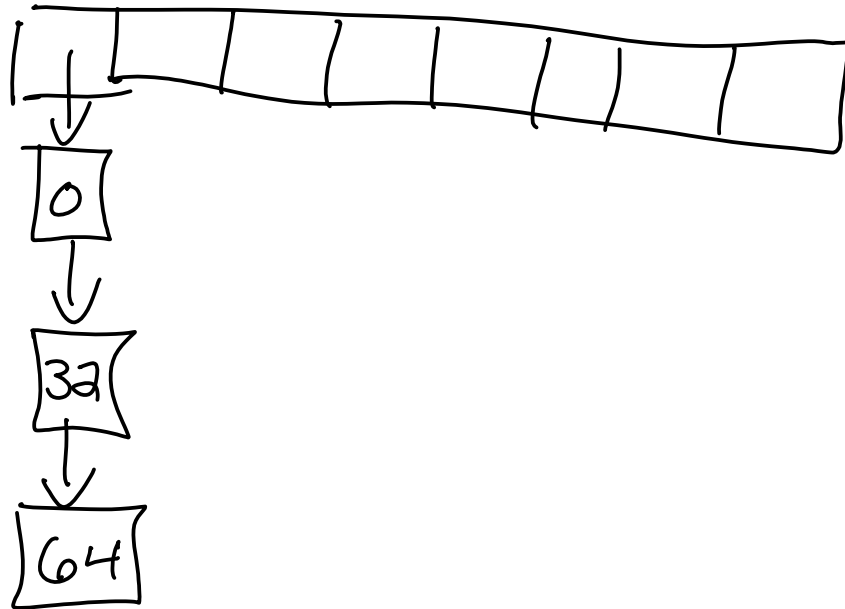
```

AB  
65 66 8 4 21  
1000001  
 11111  
 1000001  
 1000001  
 1000001  
 .....  
 .....  
 .....  
 .....  
 .....

```
int hash (int n) {  
    return n;  
}
```



insert (0, ...)  
insert (32, ...)  
insert (64, ...)  
insert (96, ...)



# Hash Table

|        |   |        |                 |
|--------|---|--------|-----------------|
| get    | { | $O(n)$ | worst good hash |
| insert |   | $O(1)$ | average         |
| update |   | $O(1)$ |                 |
| remove |   | $O(n)$ | worst bad hash  |