

### 3.1 analysis and big-O

Tuesday, September 13, 2022

Reminder: test 1 in lab next week

### TODAY

polymorphism and destructors

evaluating algorithms } empirical analysis  
theoretical analysis

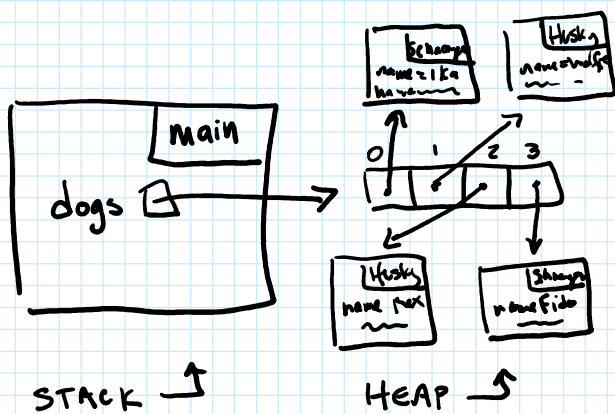
big-O definition, theoretical analysis

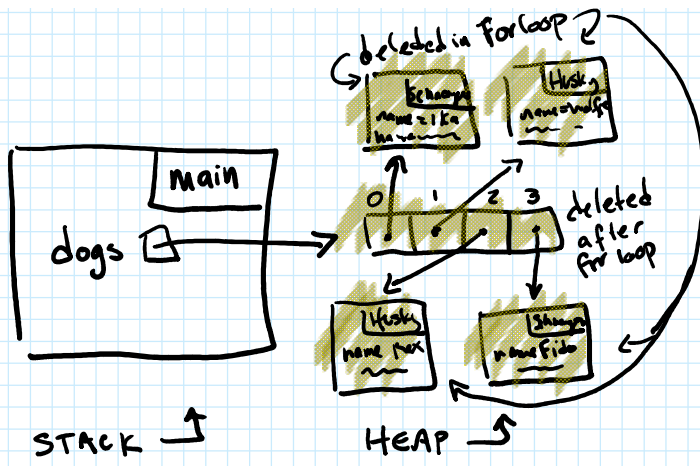
Recap: polymorphism, inheritance, destructors

With an inheritance relationship like



a pointer to the super class type Dog can point to any memory storing an object whose type is a subclass





In lab we'll see  
 valgrind ./main  
 again.

problem: bool is\_sorted(int\* array, int size)

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returns  $\left\{ \begin{array}{l} \text{true, if values are in ascending order} \\ \text{false, if values are NOT in ascending order} \end{array} \right.$

ex

int a[5] = {2, 4, 6, 8, 10};

int b[3] = {13, 1, 7};

is\_sorted(a, 5) true!

is\_sorted(b, 3) false!

Which implementation is "best"?

- fastest time to run
- smallest memory space used
- readable, easy to understand
- flexible, easy to edit
- correct

version 1: correct ✓

version 2: correct ✓

version 3: correct ✓

## EMPIRICAL ANALYSIS

Want to compare versions to find which is most efficient.

- let's run some experiments and time them
- vary the size of the array

version	input size	10	100	1000	10,000	50k
1		0.004	0.004	0.009	0.125	2.9
2		0.005	0.003	0.002	0.004	0.004
3		0.006	0.016	0.091	0.857	4.2

overall:

version 2 seems best

versions 1 & 3 are both slower,

but 3 is worse than 1 as size increases

notes on empirical analysis

- clock time is not the cleanest measure  
hard to control conditions for equal runs
- can only run on a limited set of values
- must implement each algorithm to test it

## THEORETICAL ANALYSIS

- abstract away from a particular language
- analyze how much work the algorithm requires

work = # of steps

- count the # of steps in an algorithm
- focus on key steps  
is - sorted: key step is comparison between elements
- let n represent the size of the problem

### Version 1

size = n

i loop goes from  $i = 0$  to  $i = \text{size} - 1$

j loop goes from  $j = i + 1$  to  $j = \text{size} - 1$   
comparison

How many comparisons are done in total?

$$(n-1) + (n-2) + (n-3) + \dots + 3 + 2 + 1 = \text{total}$$
$$1 + 2 + 3 + \dots + (n-3) + (n-2) + (n-1) = \text{total}$$

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$$n = n + n + \dots + n + n + n = 2 \cdot \text{total}$$

$$(n)(n-1) = 2 \cdot \text{total}$$

$$\text{total} = \frac{n^2}{2} - \frac{n}{2}$$

### Version 2:

Version 2:

i loop goes from  $i=0$  to  $\text{size}-1$   
comparison

How many comparisons are done in total?

Version 3

i loop goes from  $i=0$  to  $\text{size}-1$   
10k comparisons

How many comparisons are done in total?