

CS 43: Computer Networks

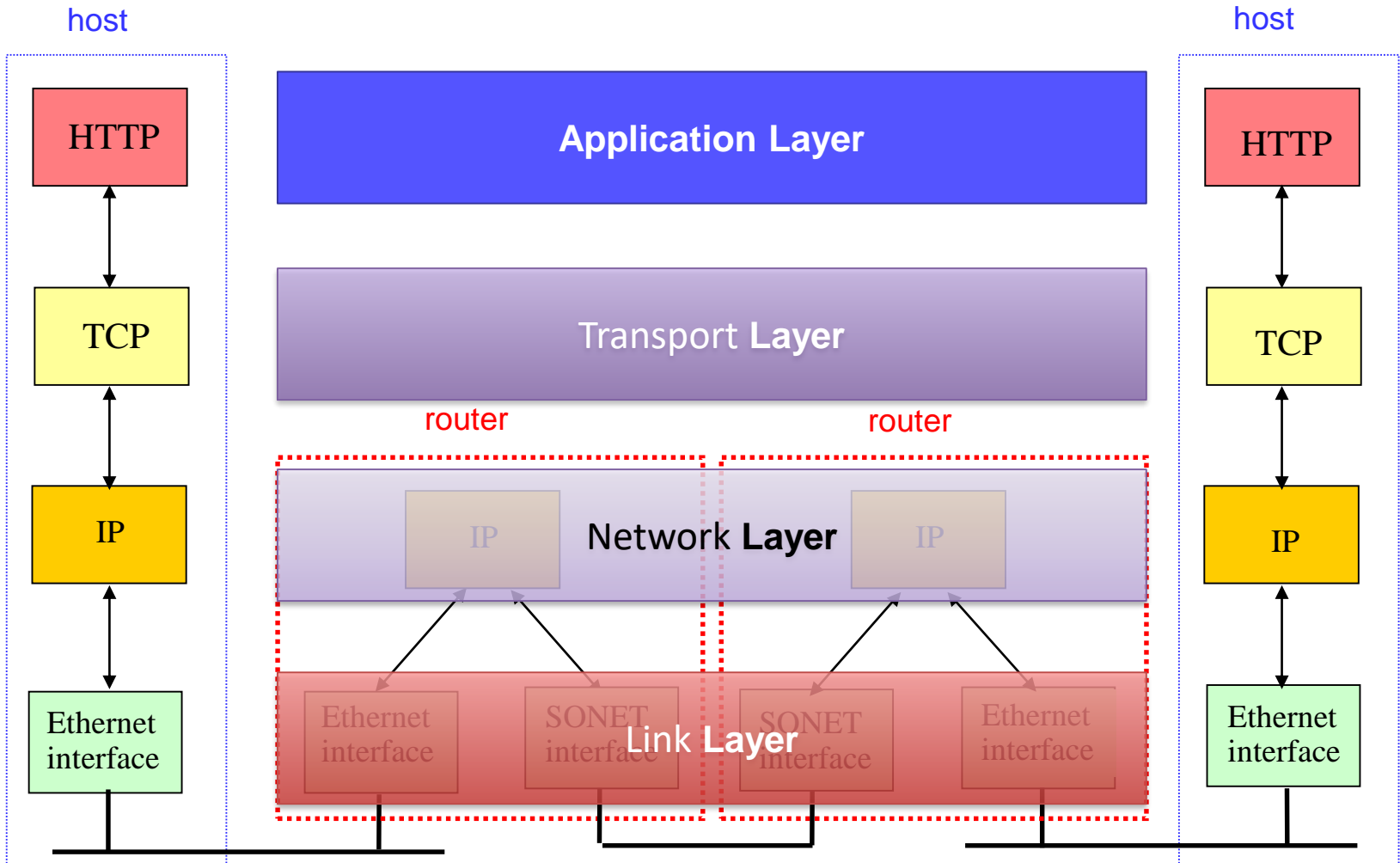
The Network Layer

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TCP/IP Protocol Stack



Some background...

- 1968: DARPAnet/ARPAnet (precursor to Internet)
 - (Defense) Advanced Research Projects Agency Network
 - Bob Taylor, Larry Roberts create program to build first wide-area packet-switched network
- Mid 1970's: new networks emerge
 - SATNet, Packet Radio, Ethernet
 - All “islands” to themselves – didn't work together
- Big question: how to connect these networks?

Internetworking

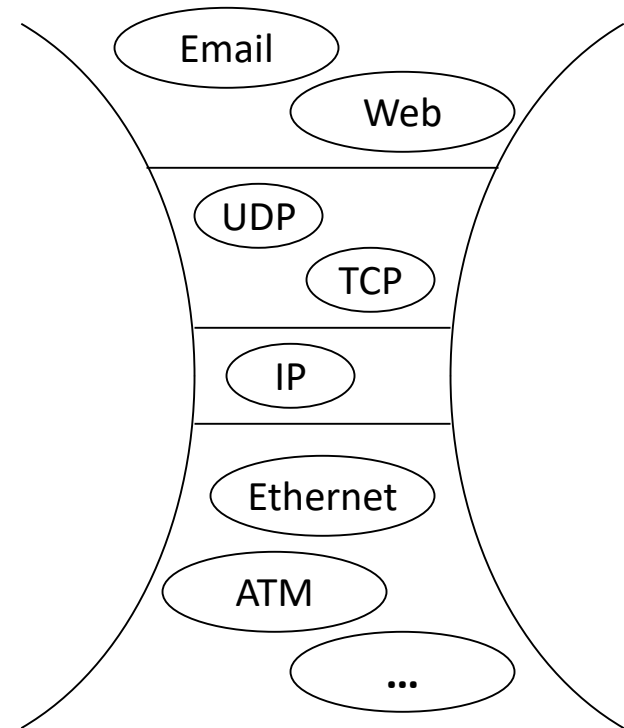
- Cerf & Kahn, in 1974,
“A Protocol for Packet Network Intercommunication”
 - Foundation for the modern Internet
- Routers forward packets from source to destination
 - May cross many separate networks along the way
- All packets use a common Internet Protocol
 - Any underlying data link protocol
 - Any higher layer transport protocol

DARPAnet Primary Goal: Connect Stuff

- “Effective technique for multiplexed utilization of existing interconnected networks” – David Clark (1988)
 - **Minimal** assumptions about underlying networks
 - No support for broadcast, multicast, real-time, reliability
 - Extra support could actually get in the way
 - Packet switched, store and forward
 - Matched application needs, nets already packet switched
 - Enables **efficient resource sharing**/high utilization
 - “Gateways” interconnect networks
 - Routers in today’s nomenclature

Internet Protocol Stack

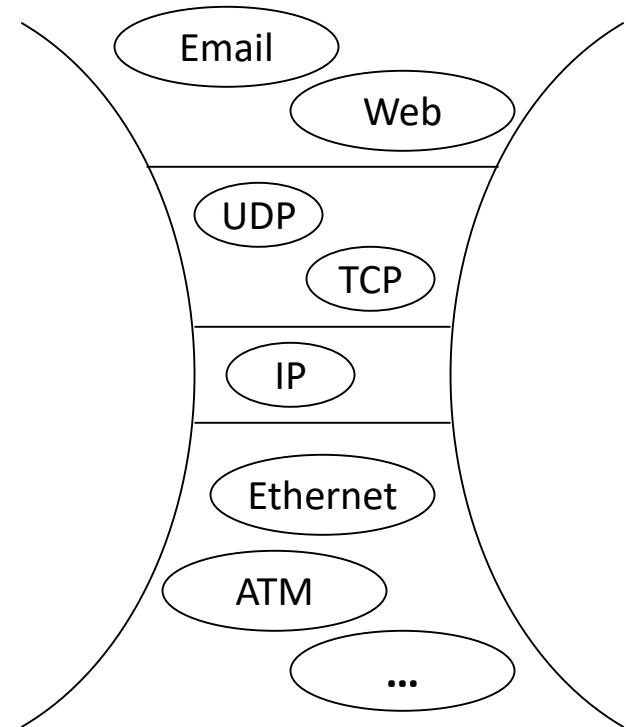
- Application: Email, Web, ...
- Transport: TCP, UDP, ...
- Network: IP
- Link: Ethernet, WiFi, ATM, ...
- Physical: copper, fiber, air, ...



- “Hourglass” model, “thin waist”, “narrow waist”

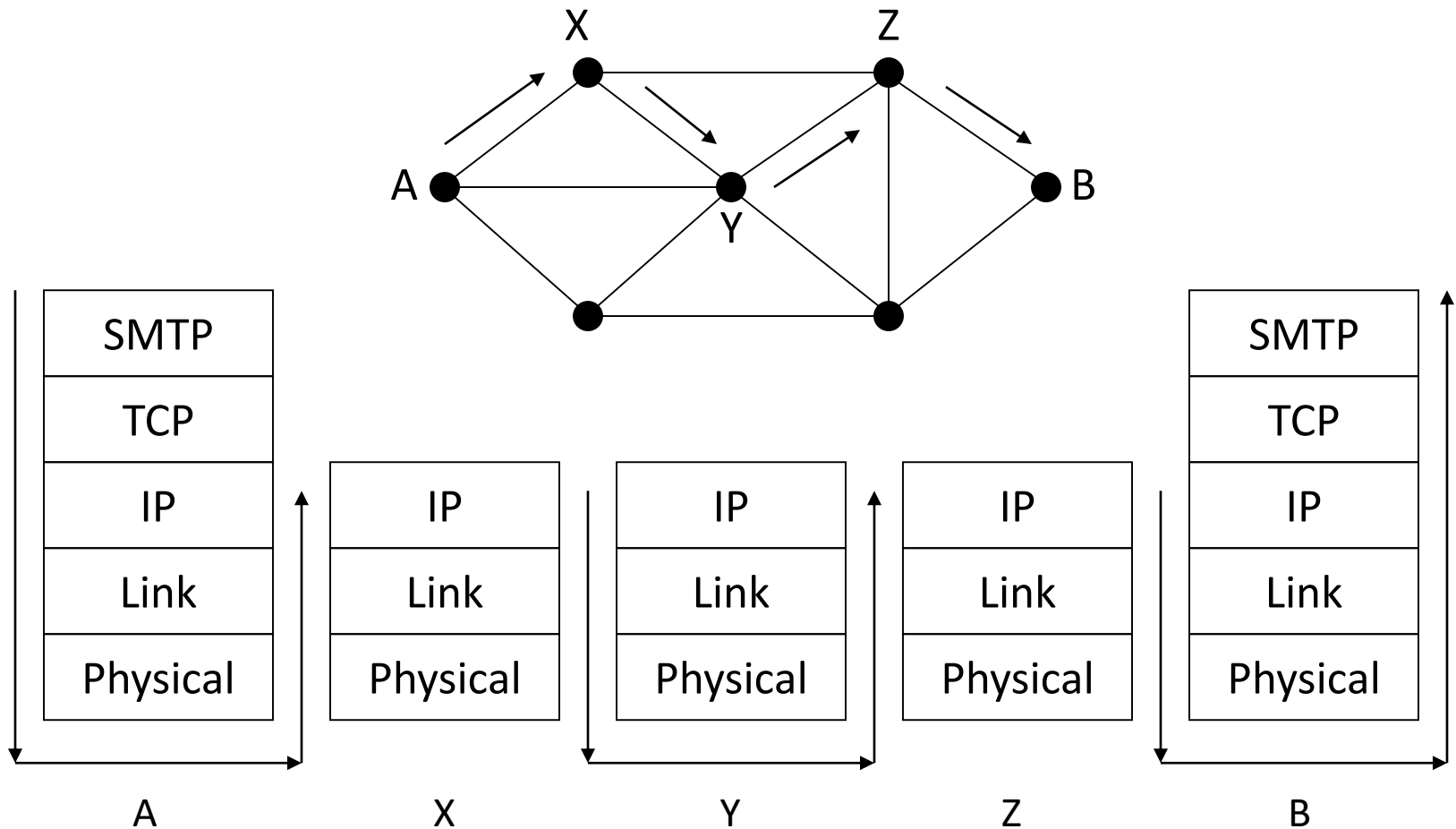
Internet Protocol Stack

- This should seem weird.
- *Everyone* uses IP?



- “Hourglass” model, “thin waist”, “narrow waist”

Example of Internet Routing



Network layer involved at every hop along the path.

Network Layer Functions

- *Forwarding*: move packets from router's input to appropriate router output ("data plane")
- *Routing*: determine route taken by packets from source to destination. ("control plane")

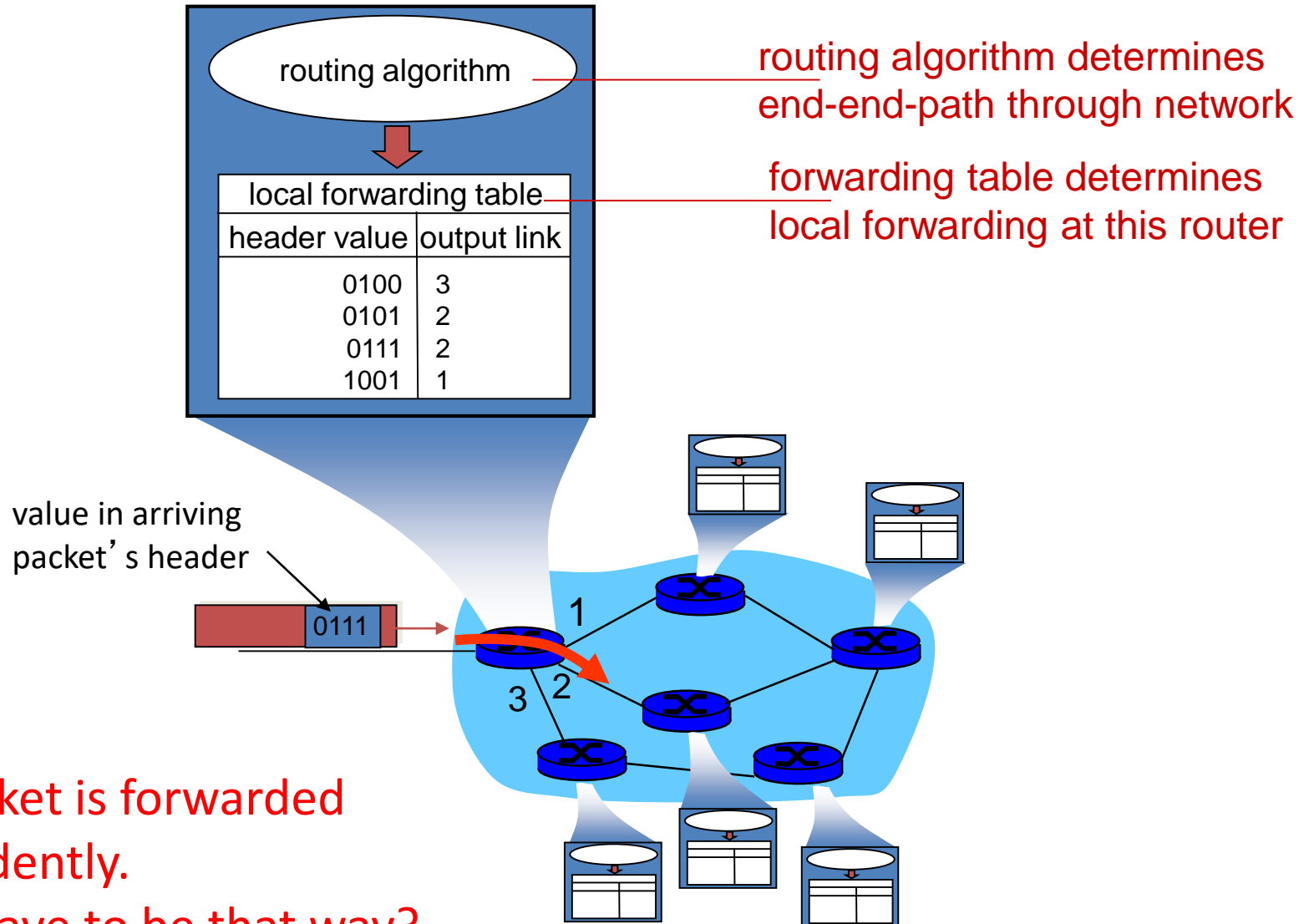
When should a router perform routing? Forwarding?

- A. Do both when a packet arrives.
- B. Route in advance, forward when a packet arrives.
- C. Forward in advance, route when a packet arrives.
- D. Do both in advance.
- E. Some other combination

Network Layer Functions

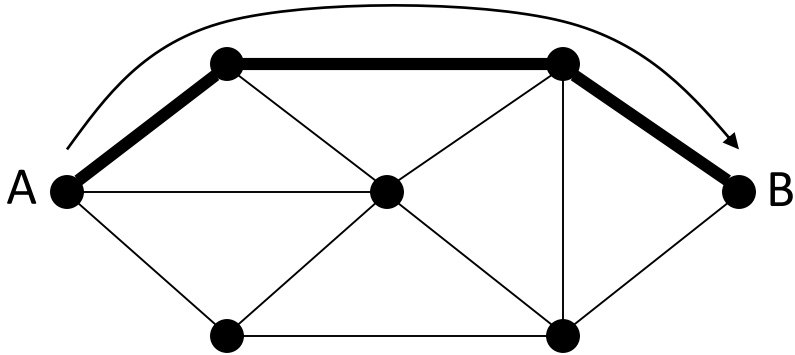
- *Forwarding*: move packets from router's input to appropriate router output
 - Look up in a table
- *Routing*: determine route taken by packets from source to destination.
 - Populating the table

Interplay between routing and forwarding



Circuit Switching

- Reserve path in advance



- (Old) telephone system

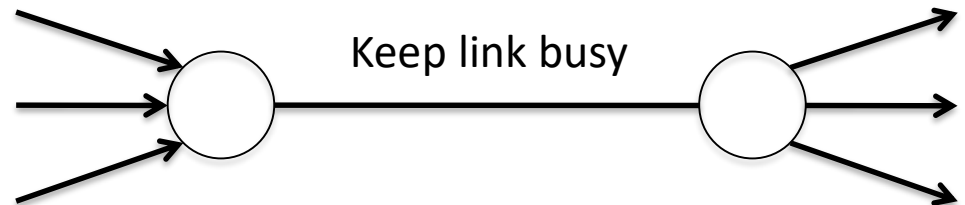
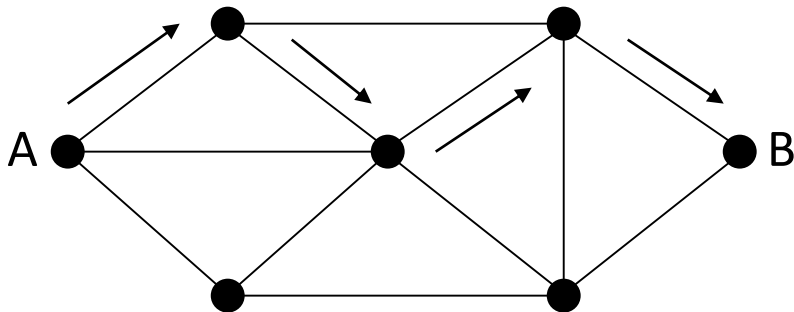


Why doesn't the Internet (typically) use circuits?

- A. It's too slow to establish a connection.
- B. It doesn't offer good enough performance.
- C. It wastes resources.
- D. It requires too many resources.
- E. Some other reason.

Packet Switching

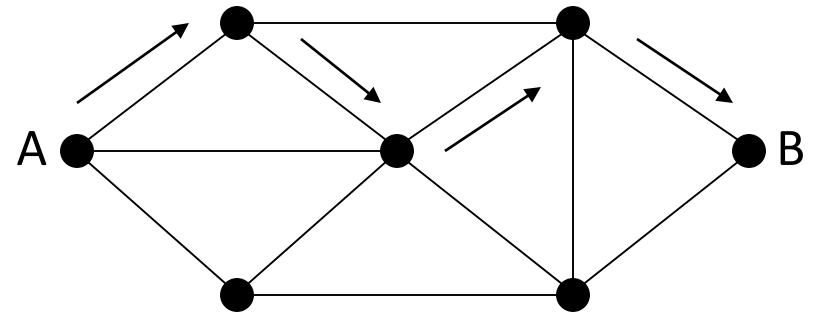
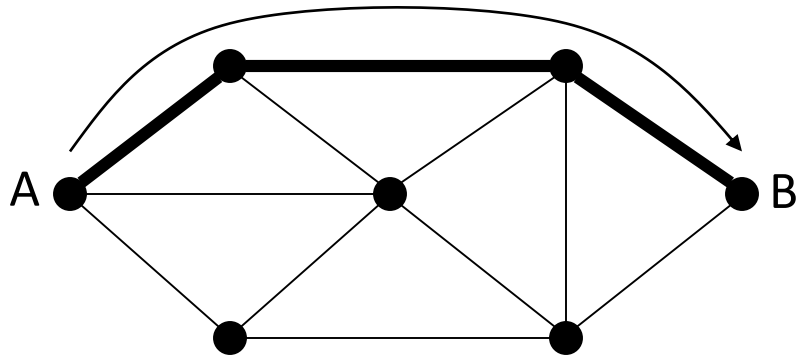
- Do we always need to reserve a link?
- Statistical multiplexing
 - Assign multiple conversations to a physical path
 - At any given time, one will have something to say



Which of the following is/are generally true of packet vs. circuit switching?

1. Packet switching has less variance in performance.
 2. Circuit switching is less reliable.
-
- A. Only 1 is true.
 - B. Only 2 is true.
 - C. Both 1 and 2 are true.
 - D. Neither 1 nor 2 are true.

Circuit-switching vs. Packet switching



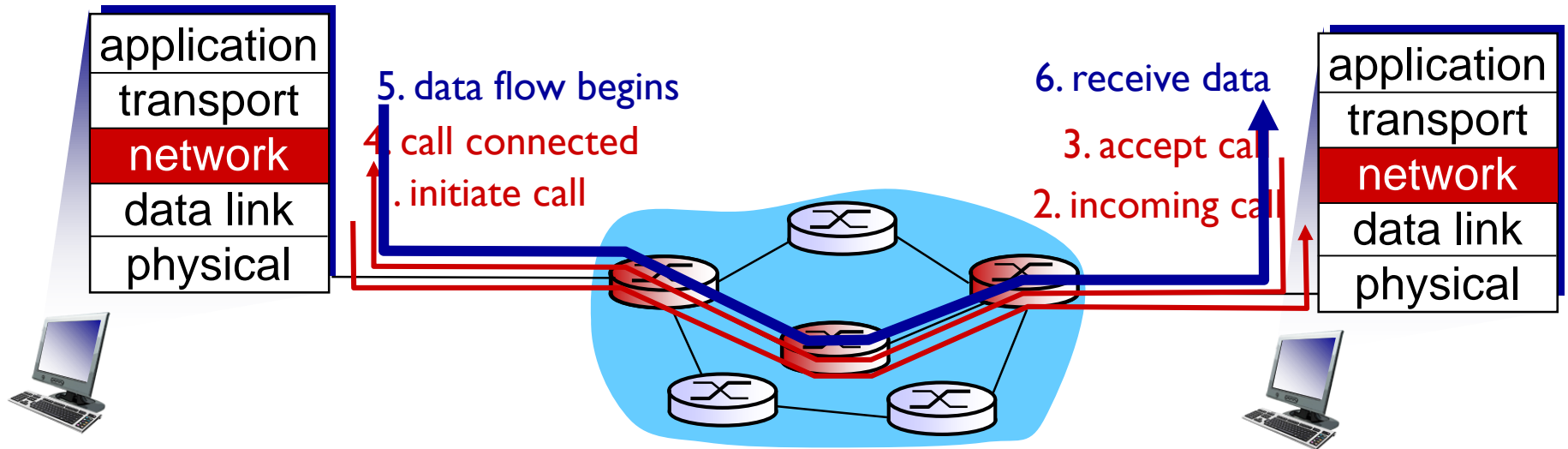
- Circuit switching: establish path, send data
 - Reserve resources, provide performance control
 - Example: telephone system
- Packet switching: forward packets hop by hop
 - Fair sharing despite bursts, statistical multiplexing
 - Example: postal system

Datagram vs. “Virtual Circuit”

- *Datagram* network provides network-layer *connectionless* service (packet switching)
- *Virtual-circuit* network provides network-layer *connection* service (like circuit switching)

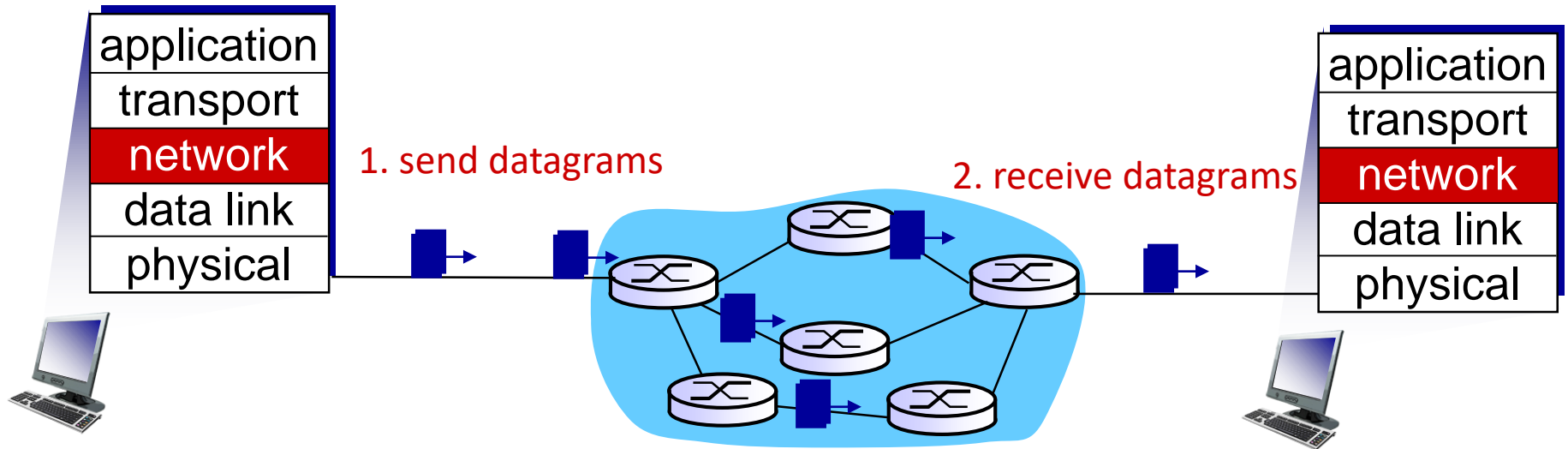
Virtual circuits: Signaling Protocols

- Used to setup, maintain, teardown VC
- Used in ATM, frame-relay, X.25
- Less common in today's Internet



Datagram Networks

- No call setup at network layer
- Routers: no state about end-to-end connections
 - no network-level concept of “connection”
- Packets forwarded individually towards destination



How do we populate a router's forwarding table?

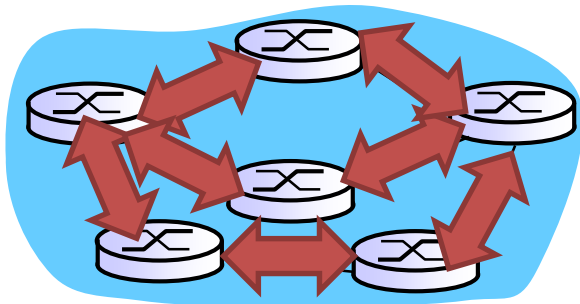
- A. A person adds entries to the table.
- B. A program external to the router adds entries to the table.
- C. Routers communicate with each other to add entries to the table.
- D. Some other mechanism.

Routing

Traditional

- Routers run a **routing protocol** to exchange state.
- Use state to build up the forwarding table.

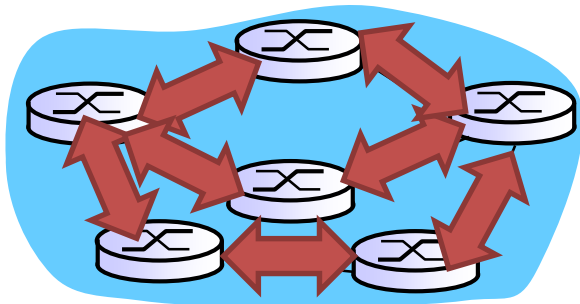
Assume this is the type of routing we're talking about unless we explicitly say otherwise!



Routing

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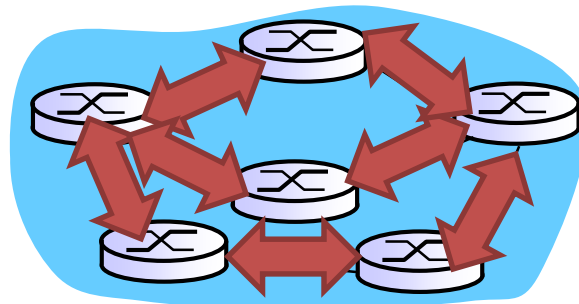


“Software-Defined”

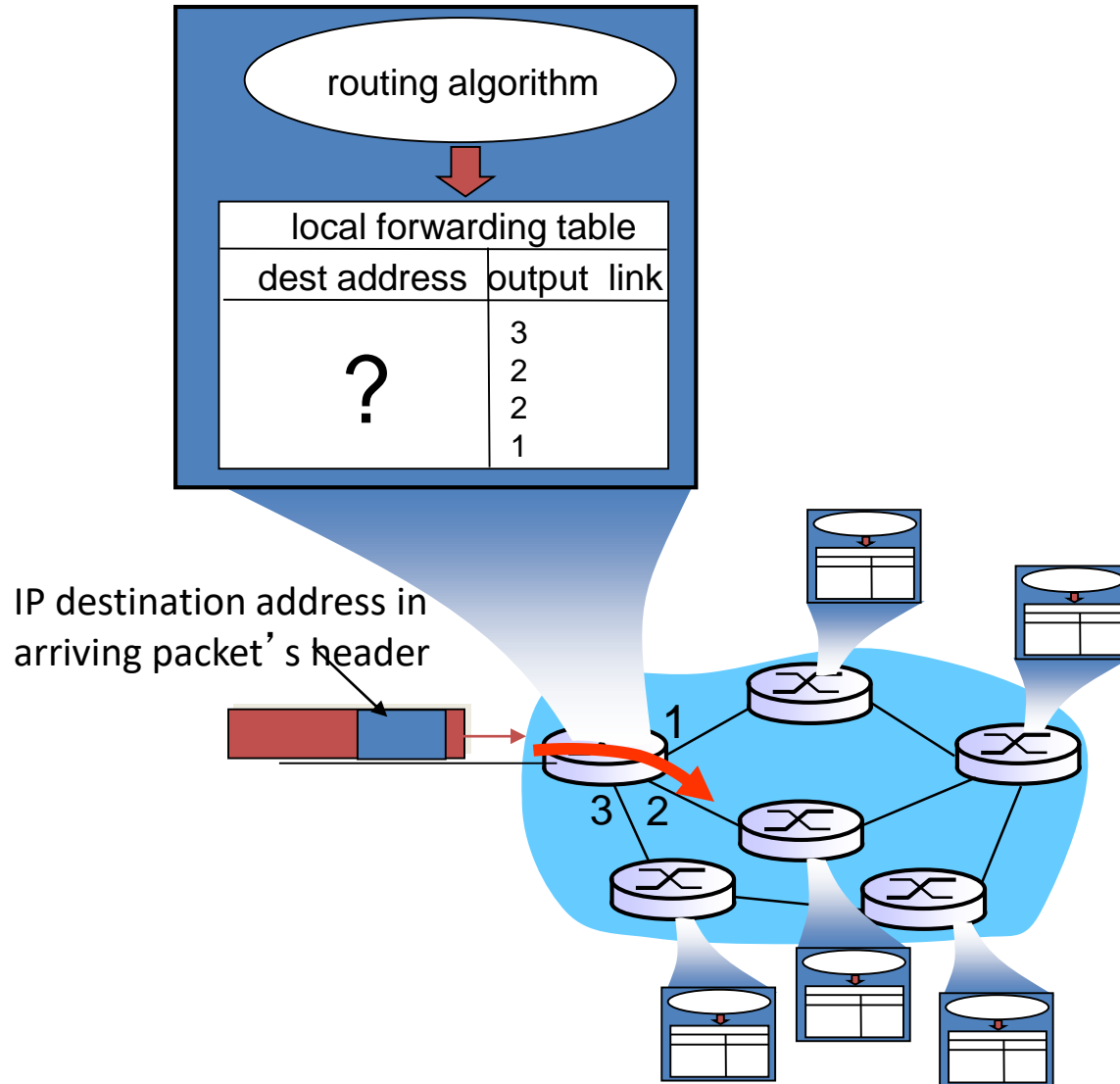
- Routers are dumb, just do what they’re told.
- Controller service explicitly tells each router what to do.
- Rare on the Internet, hot topic in data centers.

Datagram Forwarding

- Routers periodically exchange state.
- Use the state to build a **forwarding table** (FIB)



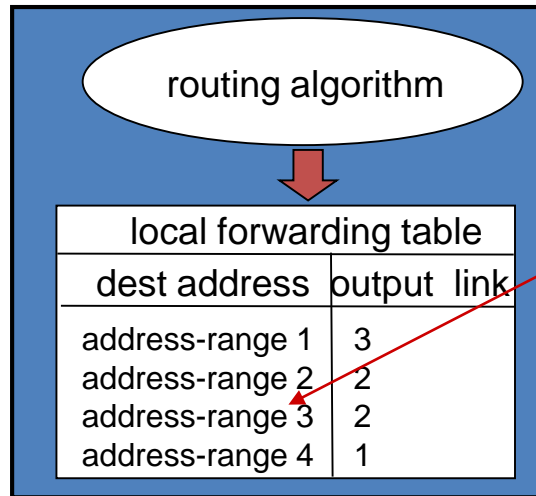
Datagram forwarding table



Routers exchange state (we'll save the what and when for later). They decide, for each destination, how to get there, and build a lookup structure. What should they build?

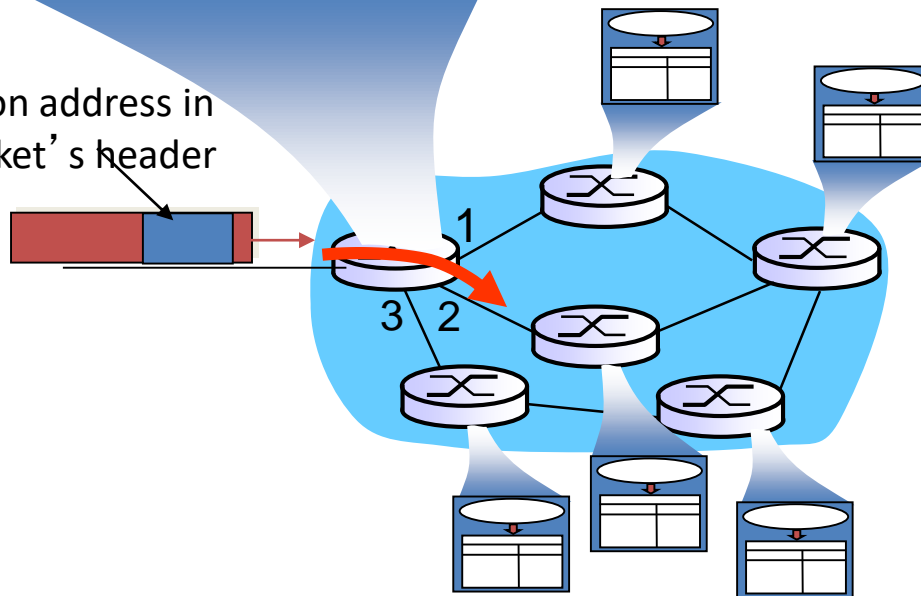
- A. A list – scan for the destination.
- B. A hash table – look up the destination.
- C. A tree – Follow branches that lead to the destination.
- D. Some other software structure.
- E. We can't do this in software, we need special hardware.

Datagram forwarding table



4 billion IP addresses, so rather than list individual destination address list *range* of addresses (aggregate table entries)

IP destination address in arriving packet's header



Datagram forwarding table

Destination Address Range	Link Interface
11001000 00010111 00010000 00000000 through 11001000 00010111 00010111 11111111	0
11001000 00010111 00011000 00000000 through 11001000 00010111 00011000 11111111	1
11001000 00010111 00011001 00000000 through 11001000 00010111 00011111 11111111	2
Otherwise (default gateway)	3

Q: but what happens if ranges don't divide up so nicely?

Longest prefix matching

longest prefix matching

when looking for forwarding table entry for given destination address, use *longest* address prefix that matches destination address.

Destination Address Range	Link interface
11001000 00010111 00010*** *****	0
11001000 00010111 00011000 *****	1
11001000 00010111 00011*** *****	2
Otherwise (default gateway)	3

examples:

DA: 11001000 00010111 00010110 10100001

which interface?

DA: 11001000 00010111 00011000 10101010

which interface?

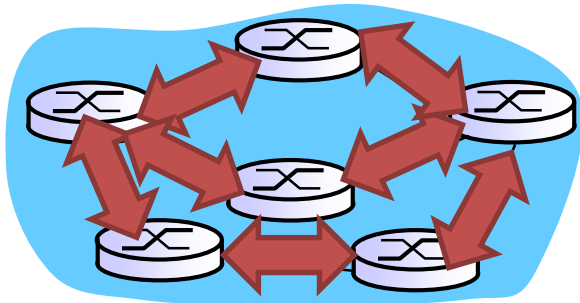
Coming up in ~1 week.



Routing

Traditional

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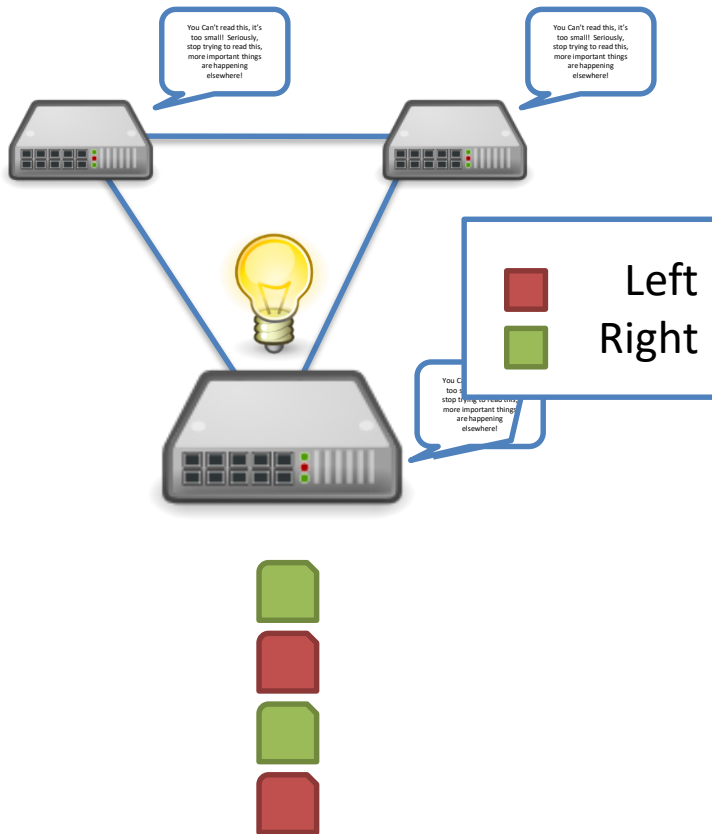


“Software-Defined”

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Software-Defined Networking (SDN)

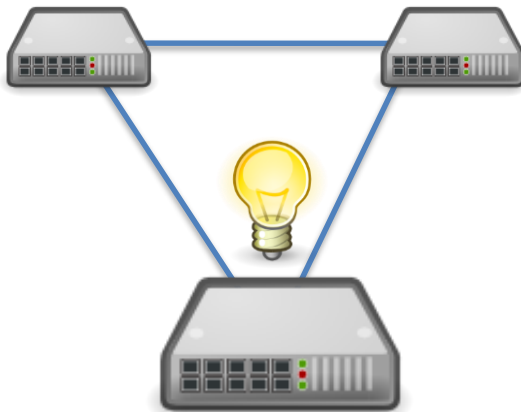
Traditional Hardware



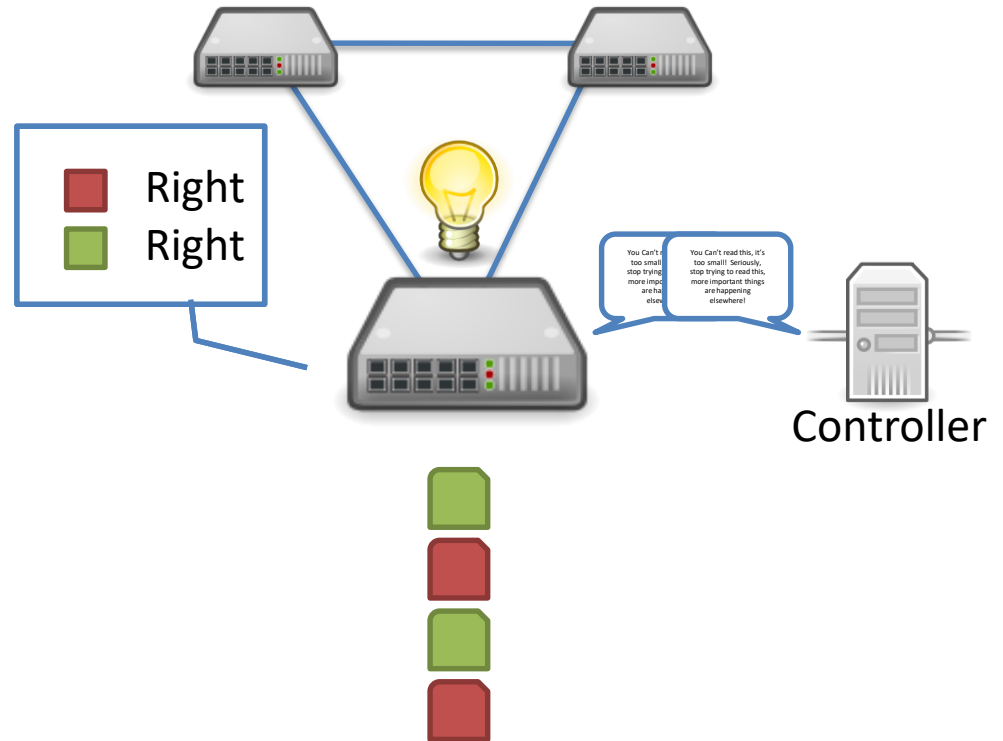
SDN Hardware

Software-Defined Networking (SDN)

Traditional Hardware



SDN Hardware



Summary

- Forwarding: moving packet from one interface to another (table lookup)
- Routing: Populating the table in advance
- On the Internet, best effort packet switching is the norm
- Hardware helps with quick forwarding using longest prefix matching