

# CS 43: Computer Networks

## Naming and DNS

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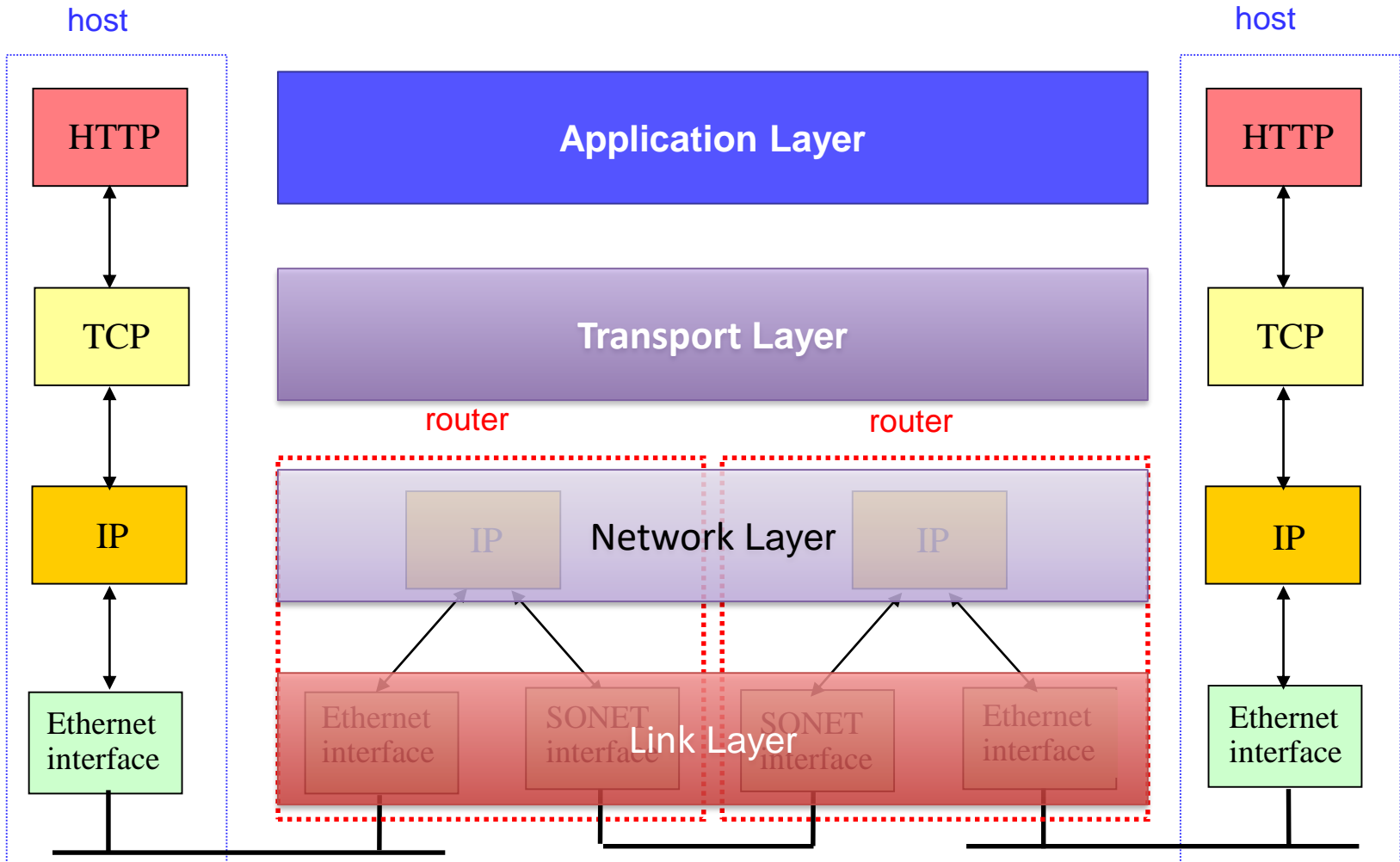
Swarthmore College

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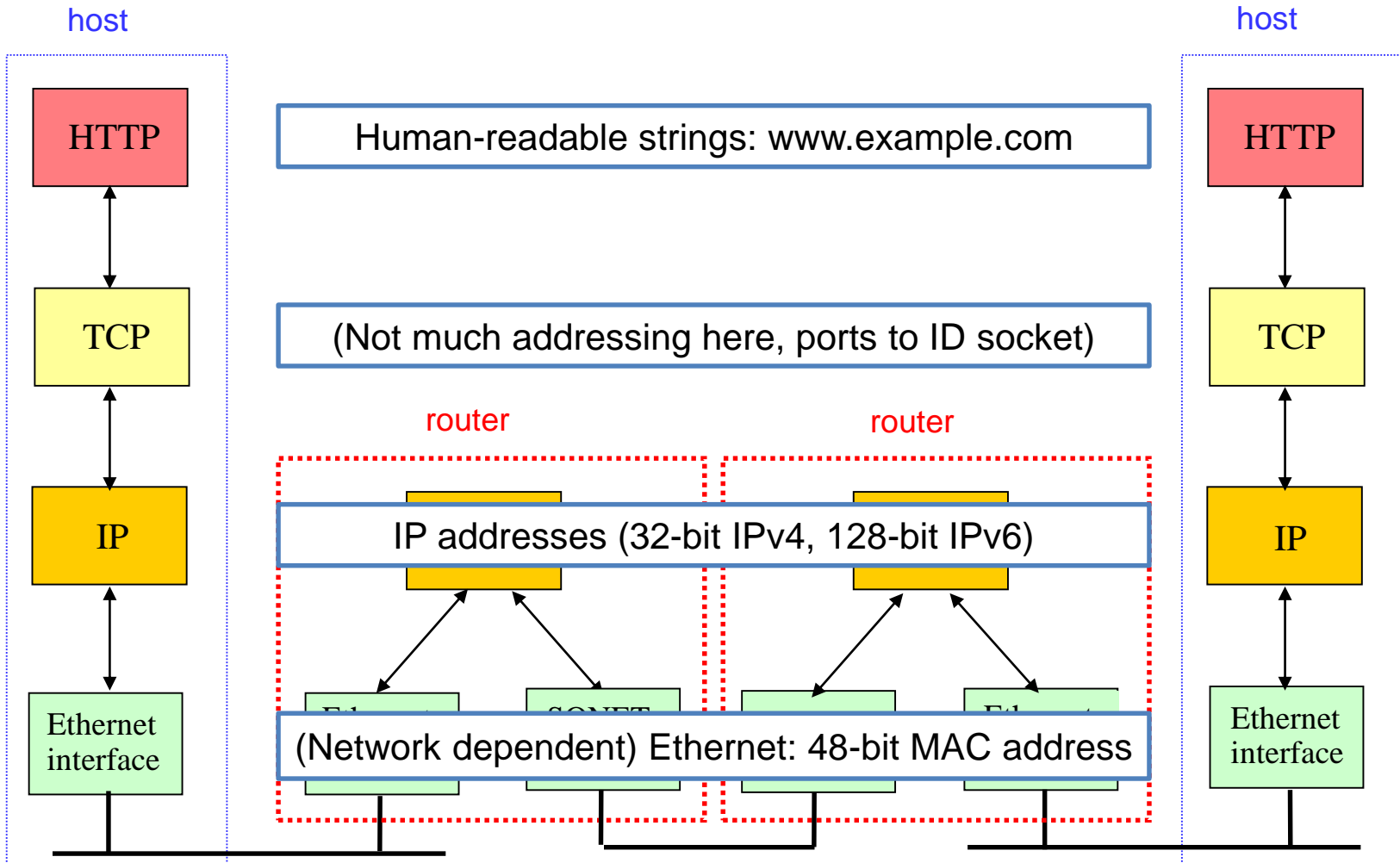
# Agenda

- Identifiers and addressing
- Domain Name System
  - History
  - Query sequences
  - Record types
  - Load balancing

# Recall: TCP/IP Protocol Stack



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# Identifiers

- **Host name** (e.g., `www.swarthmore.edu`)
  - Used by *humans* to specify host of interest
  - Unique, selected by host administrator
  - Hierarchical, variable-length string of alphanumeric characters
- **IP address** (e.g., `130.58.68.164`)
  - Used by *routers* to forward packets
  - Unique, topologically meaningful locator
  - Hierarchical namespace of 32 bits
- **MAC address** (e.g., `D8:D3:85:94:5F:1E`)
  - Used by *network adaptors* to identify interesting frames
  - Unique, hard-coded identifier burned into network adaptor
  - Flat name space (of 48 bits in Ethernet)

# What's in a name?

- Host name: `web.cs.swarthmore.edu`
  - **Domain**: registrar for each top-level domain (e.g., .edu)
  - **Host name**: local administrator assigns to each host
- IP addresses: `130.58.68.164`
  - **Prefixes**: ICANN, regional Internet registries, and ISPs
  - **Hosts**: static configuration, or dynamic using DHCP
- MAC addresses: `D8:D3:85:94:5F:1E`
  - **OIDs**: assigned to vendors by the IEEE
  - **Adapters**: assigned by the vendor from its block

# What's in a name?

- Host name: `web.cs.swarthmore.edu` (today)
  - **Domain**: registrar for each top-level domain (e.g., .edu)
  - **Host name**: local administrator assigns to each host
- IP addresses: `130.58.68.164` (a few weeks)
  - **Prefixes**: ICANN, regional Internet registries, and ISPs
  - **Hosts**: static configuration, or dynamic using DHCP
- MAC addresses: `D8:D3:85:94:5F:1E`
  - **OIDs**: assigned to vendors by the IEEE
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# Mapping Between Identifiers

- Domain Name System (**DNS**)
  - Given a host name, provide the IP address
  - Given an IP address, provide the host name
- Address Resolution Protocol (**ARP**)
  - Given an IP address, provide the MAC address
  - To enable communication within the Local Area Network
- Dynamic Host Configuration Protocol (**DHCP**)
  - Automates host boot-up process
  - Given a MAC address, assign a unique IP address
  - ... and tell host other stuff about the Local Area Network



What's the biggest challenge for DNS?

# What's the biggest challenge for DNS?

- A. It's old.
- B. The fact that the Internet is global.
- C. The fact that DNS is now critical infrastructure.
- D. The sheer number of name lookups happening at any given time.
- E. How and when the name -> IP address mapping should change.

# In the old days...

- Pre-1982, everyone downloads a “hosts.txt” file from SRI
- Pre-1998, Jon Postel, researcher at USC, runs the Internet Assigned Numbers Authority (IANA)
  - RFCs 882 & 883 in 1983
  - RFCs 1034 & 1035 in 1987



- Emailed 8/12 root DNS servers, asked change to his authority. They did.
- <http://www.wired.com/wiredenterprise/2012/10/joe-postel/>

# Since 1998...

- Control of Internet Assigned Numbers Authority (IANA) transferred to Internet Corporation for Assigned Names and Numbers (ICANN)
  - ICANN is a private non-profit (formerly) blessed by US DOC
  - Global advisory committee for dealing with international issues
  - 2000's: Many efforts for UN control, US resisted
  - 2016: ICANN no longer partnered with DOC
- Lots of geopolitics here...

# Who should control DNS?

- A. US government
- B. UN / International government
- C. Private corporation
- D. Someone else

# Recent Controversy

- Is ICANN working in the world's best interest?
- New “top level domains” added, for auction
- Example: the “.sucks” TLD (+ many others)

# Reality

- As computer scientists, it's probably not up to us to decide. 😞
- Let's focus on the technical aspects of DNS. 😊

# DNS Services

- DNS is an application-layer protocol. (E2E design!)
- It provides:
  - Hostname to IP address translation
  - Host aliasing (canonical and alias names)
  - Mail server aliasing
  - Load distribution (one name may resolve to multiple IP addresses)
  - Lots of other stuff that you might use a directory service to find. (Wikipedia: List of DNS record types)



# DNS Records

*DNS*: distributed DB storing resource records (**RR**)

RR format: (name, value, type, ttl)

## type=A

- **name** is hostname
- **value** is IP address

## type=NS

- **name** is domain (e.g., foo.com)
- **value** is hostname of authoritative name server for this domain

## type=CNAME

- **name** is alias name for some “canonical” (the real) name
- **www.ibm.com** is really **servereast.backup2.ibm.com**
- **value** is canonical name

## type=MX

- **value** is name of mailserver associated with **name**

# DNS protocol, messages

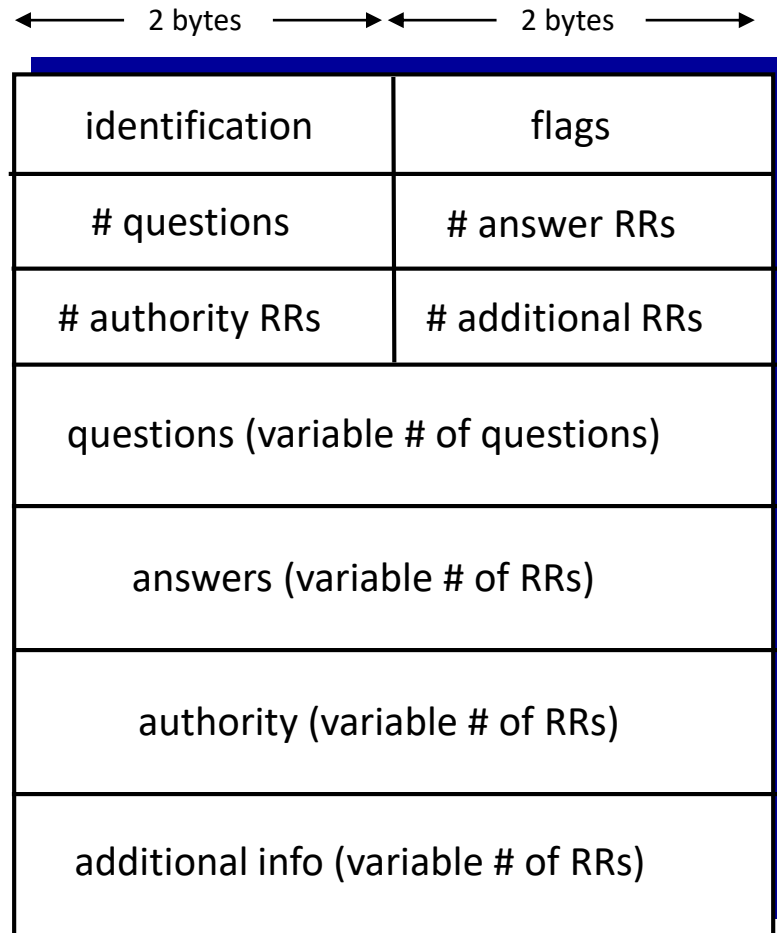
- *query* and *reply* messages, both with same *message format*

## Message header

- **identification**: 16 bit # for query, reply to query uses same #
- **flags**:
  - query or reply
  - recursion desired
  - recursion available
  - reply is authoritative

## Sent via UDP

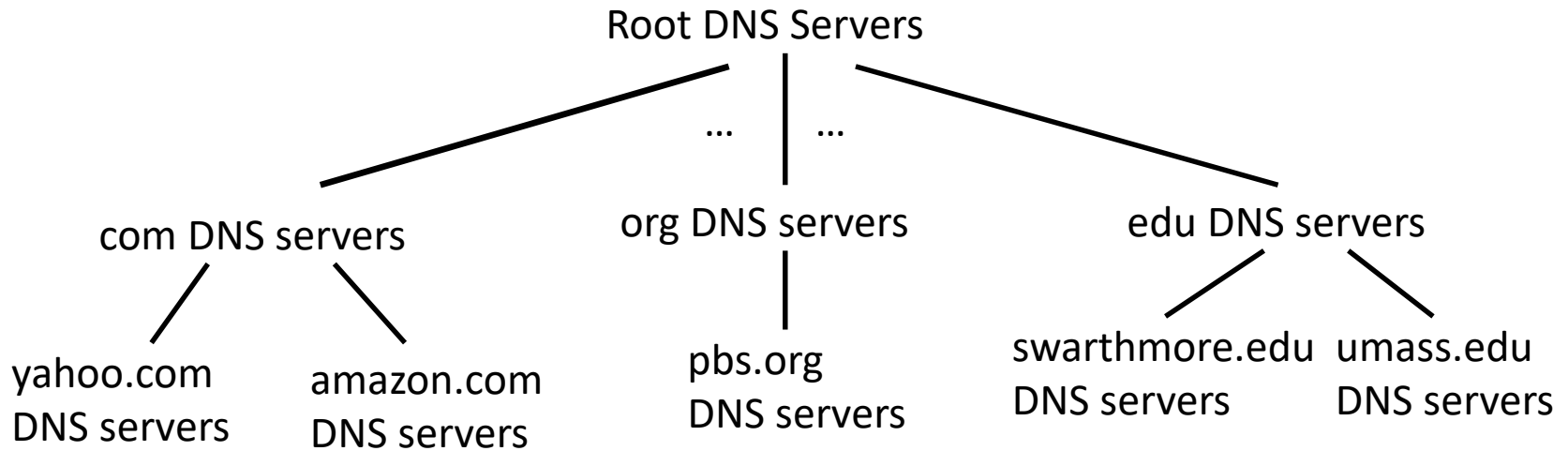
- No connection established
- Not reliable



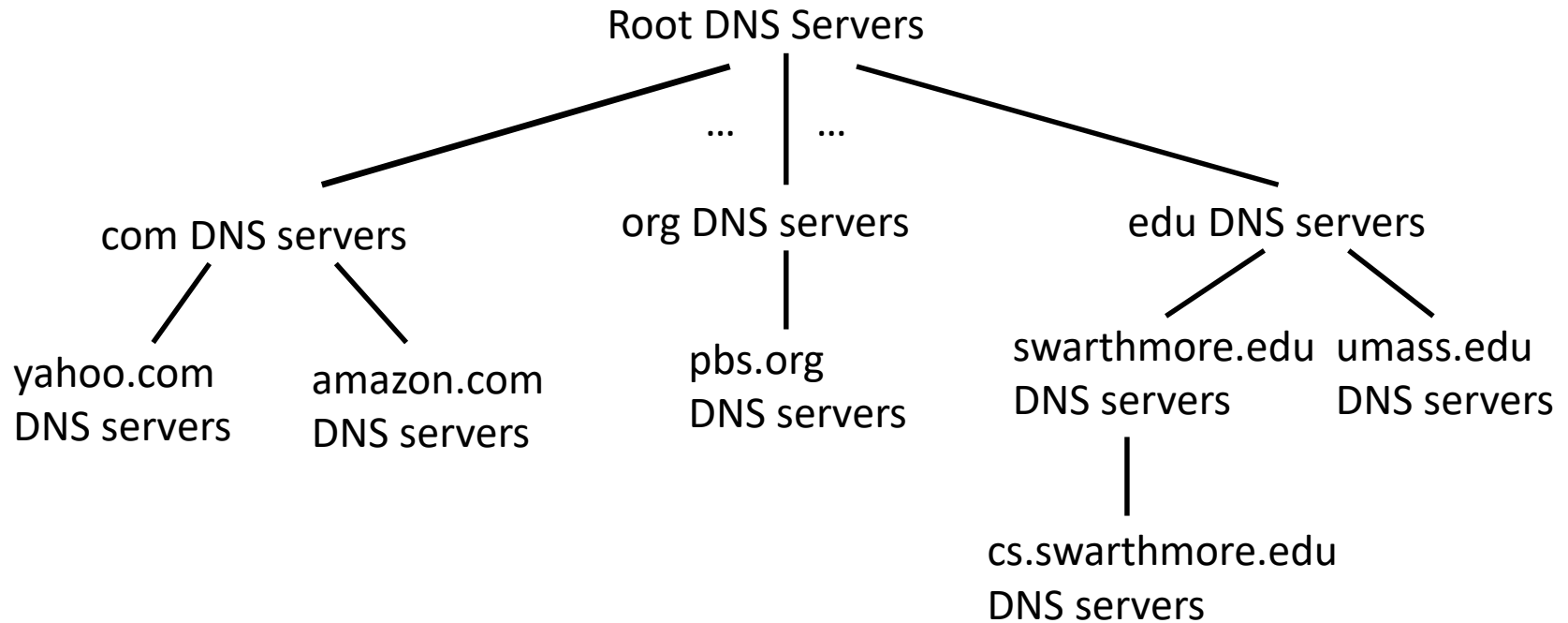
# Domain Name System (DNS)

- Distributed administrative control
  - Hierarchical name space divided into zones
  - Distributed over a collection of DNS servers
- Hierarchy of DNS servers
  - Root servers
  - Top-level domain (TLD) servers
  - Authoritative DNS servers
- Performing the translations
  - Local DNS servers
  - Resolver software

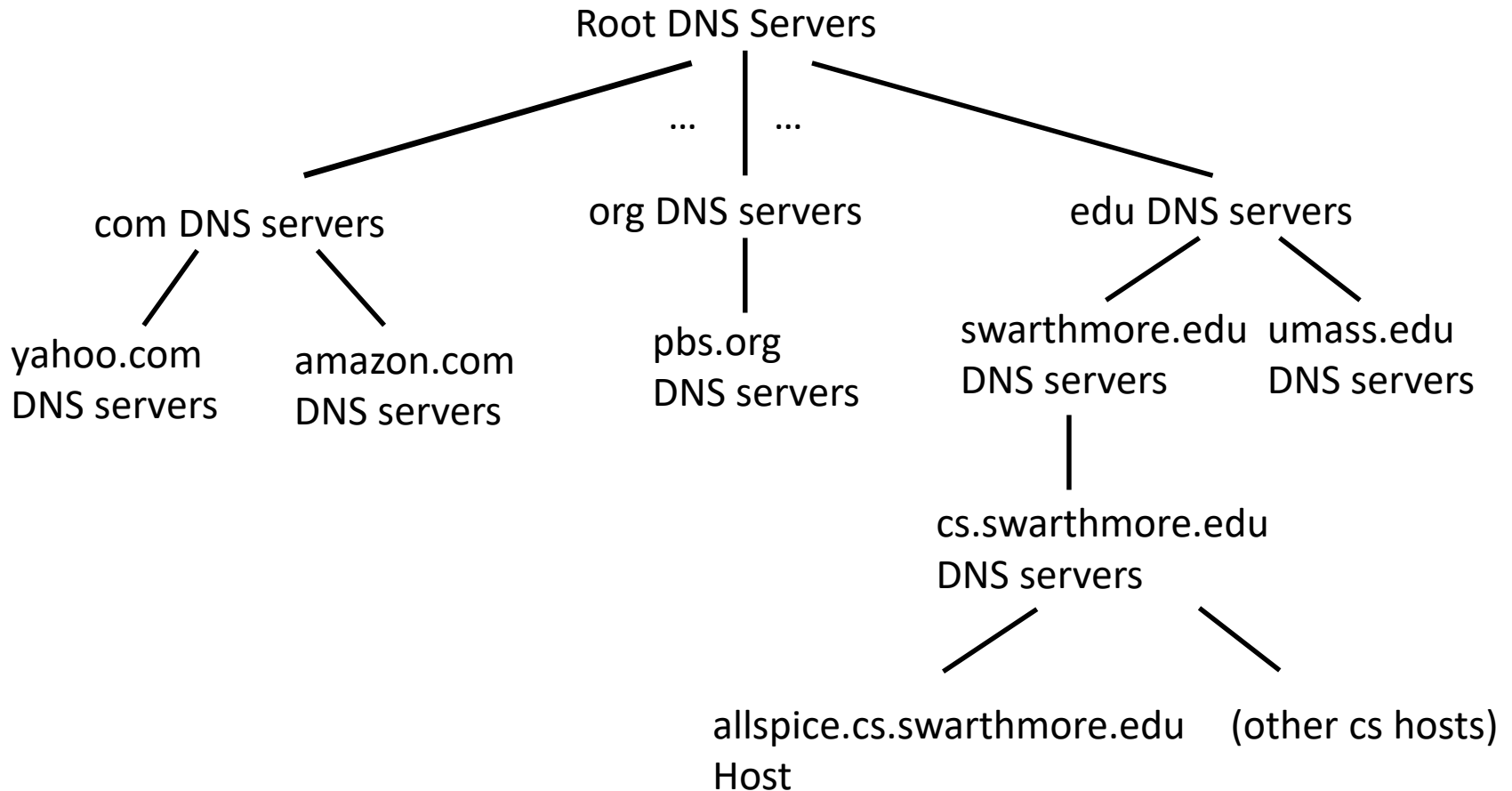
# DNS: a distributed, hierarchical database



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- allspice.cs.swarthmore.edu.

Nameless root,  
Usually implied.

# Why do we structure DNS like this?

## Which of these helps the most?

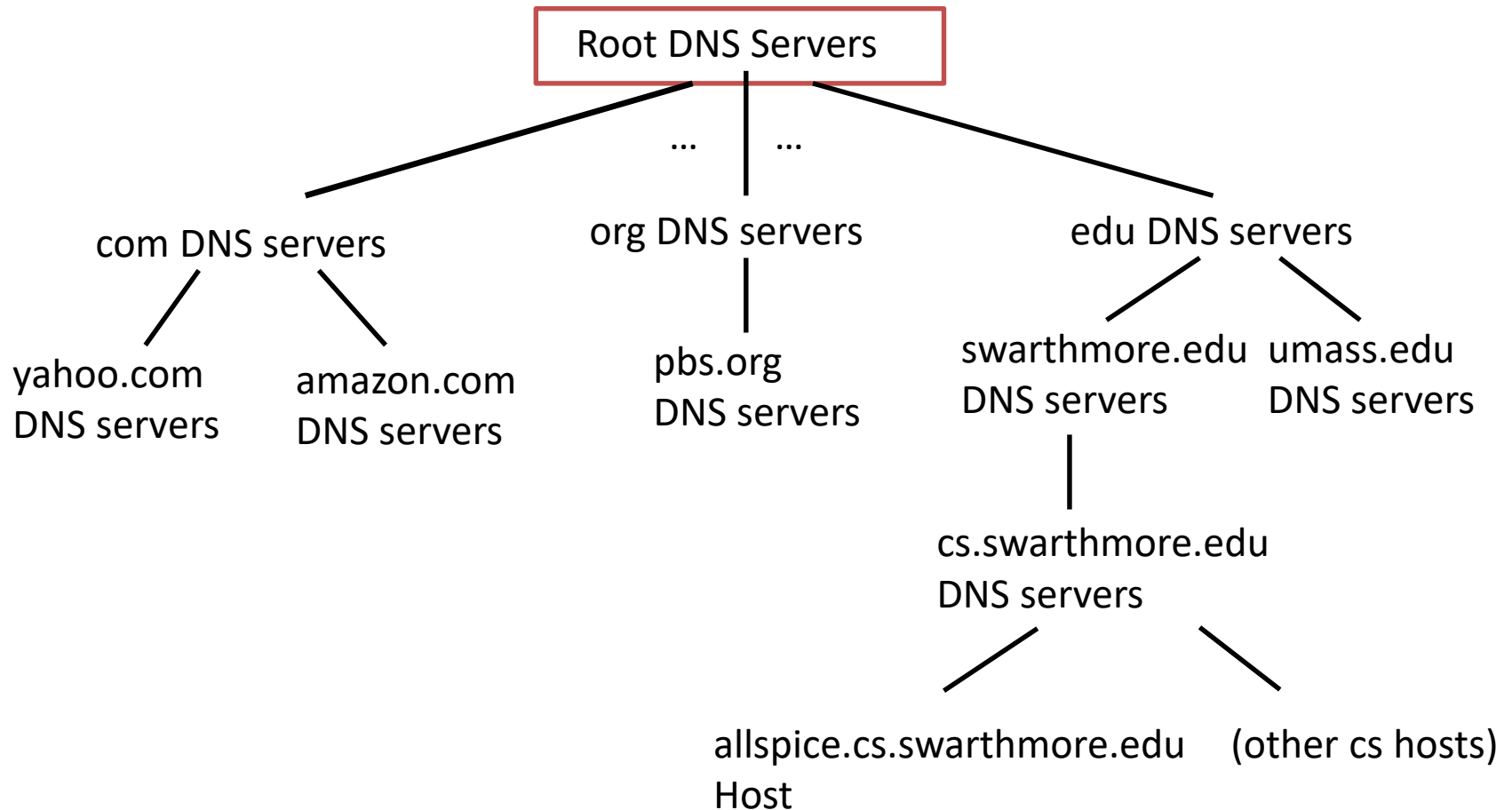
A. It divides up responsibility among parties.

B. It improves performance of the system.

C. It reduces the size of the state that a server needs to store.

D. Some other reason.

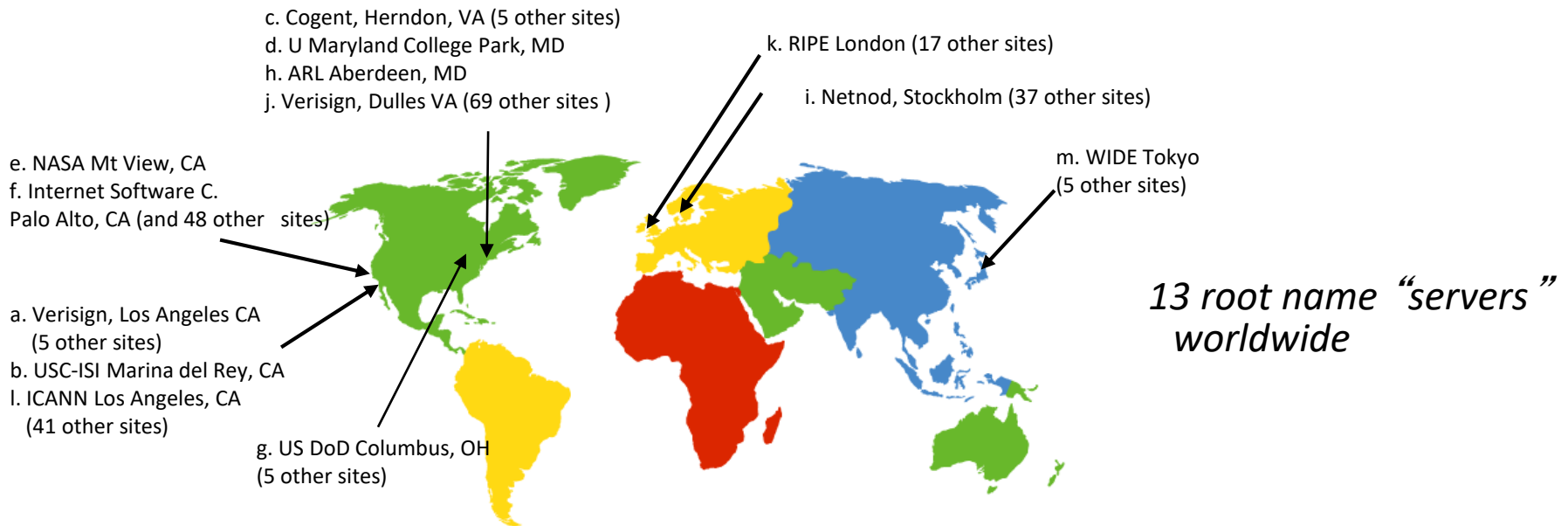
# DNS: a distributed, hierarchical database





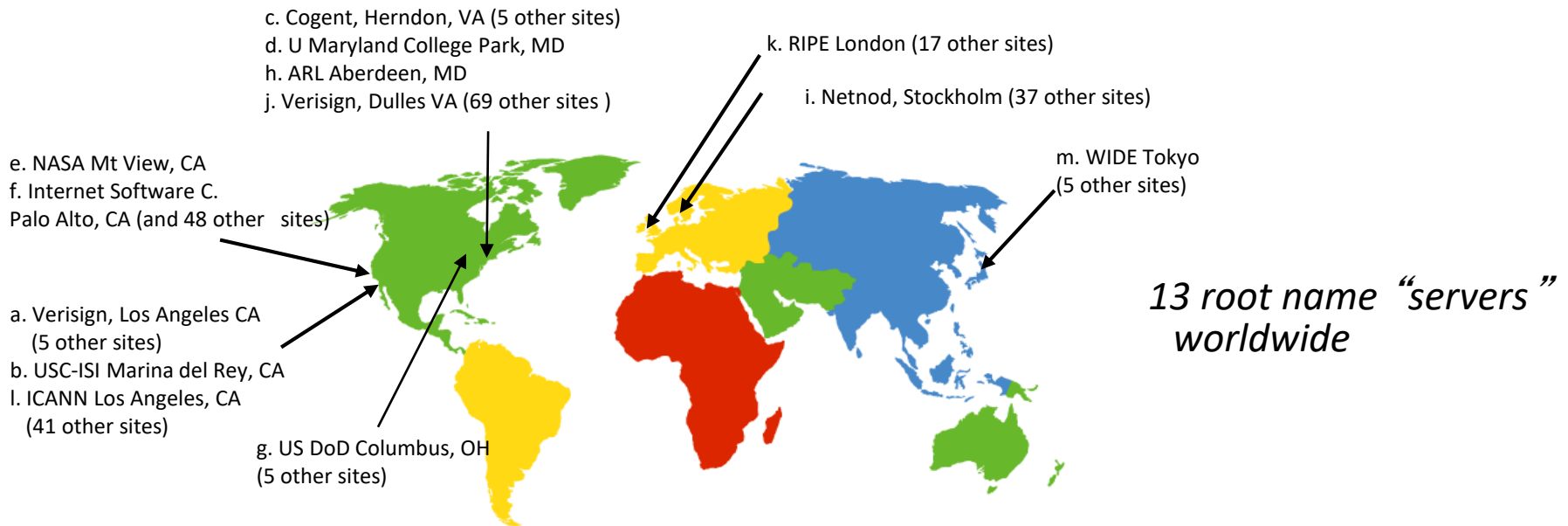
# DNS: Root Name Servers

- Root name server:
  - Knows how to find top-level domains (.com, .edu, .gov, etc.)
  - How often does the location of a TLD change?

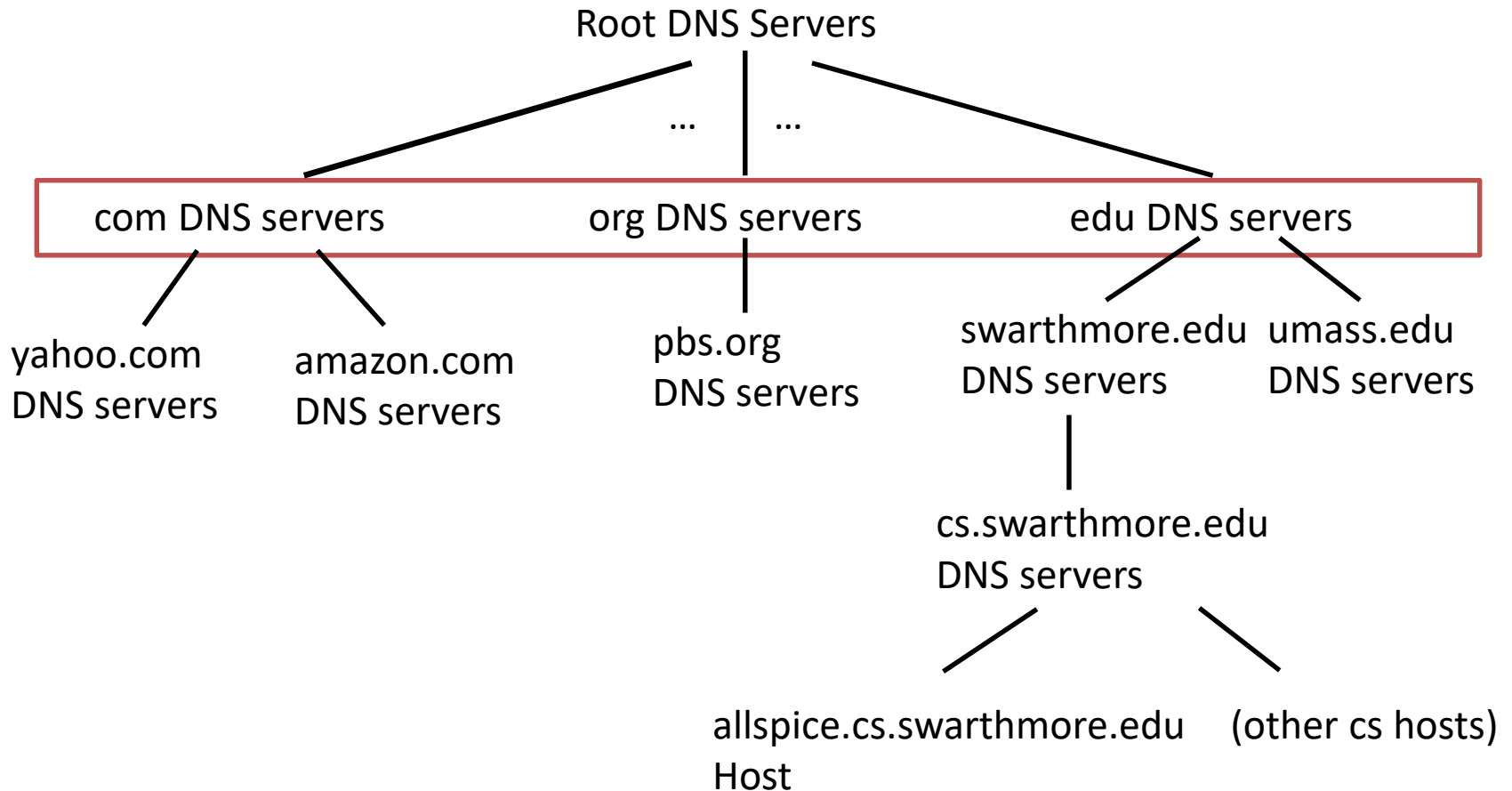


# DNS: Root Name Servers

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  - How often does the location of a TLD change?
  - ~300 total root servers
  - Significant amount of traffic is not legitimate



# DNS: a distributed, hierarchical database

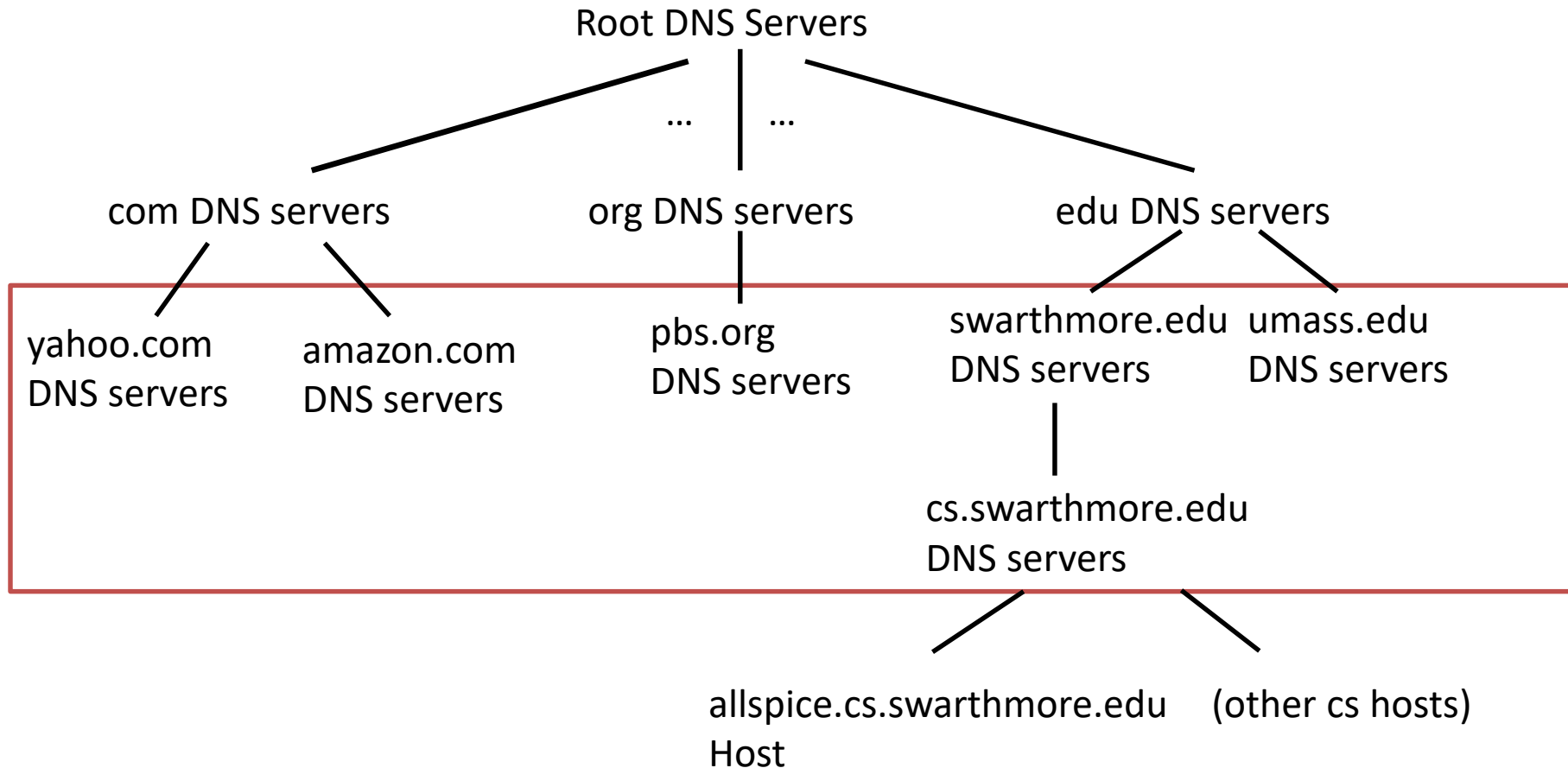


# Top Level Domains

## *Top-level domain (TLD) servers:*

- Responsible for com, org, net, edu, gov, aero, jobs, museums, and all top-level country domains, e.g.: uk, fr, de, ca, jp, etc.
- Verisign maintains servers for .com and .net TLD
- Educause for .edu TLD (Verisign actually runs backend)
- Others managed by corresponding entity (e.g., local governments or companies)

# DNS: a distributed, hierarchical database



# Authoritative Servers

## *Authoritative DNS servers:*

- Organization's own DNS server(s), providing authoritative hostname to IP mappings for organization's named hosts
- Can be maintained by organization or service provider, easily changing entries
- Often, but not always, acts as organization's local name server (for responding to look-ups)

# Resolution Process

- End host wants to look up a name, who should it contact?
  - It could traverse the hierarchy, starting at a root
  - More efficient for ISP to provide a local server
- ISP's local server for handling queries not necessarily a part of the pictured hierarchy

# Local DNS Name Server

- Each ISP (residential ISP, company, university) has (at least) one
  - also called “default name server”
- When host makes DNS query, query is sent to its local DNS server
  - has local cache of recent name-to-address translation pairs (but may be out of date!)
  - acts as proxy, forwards query into hierarchy

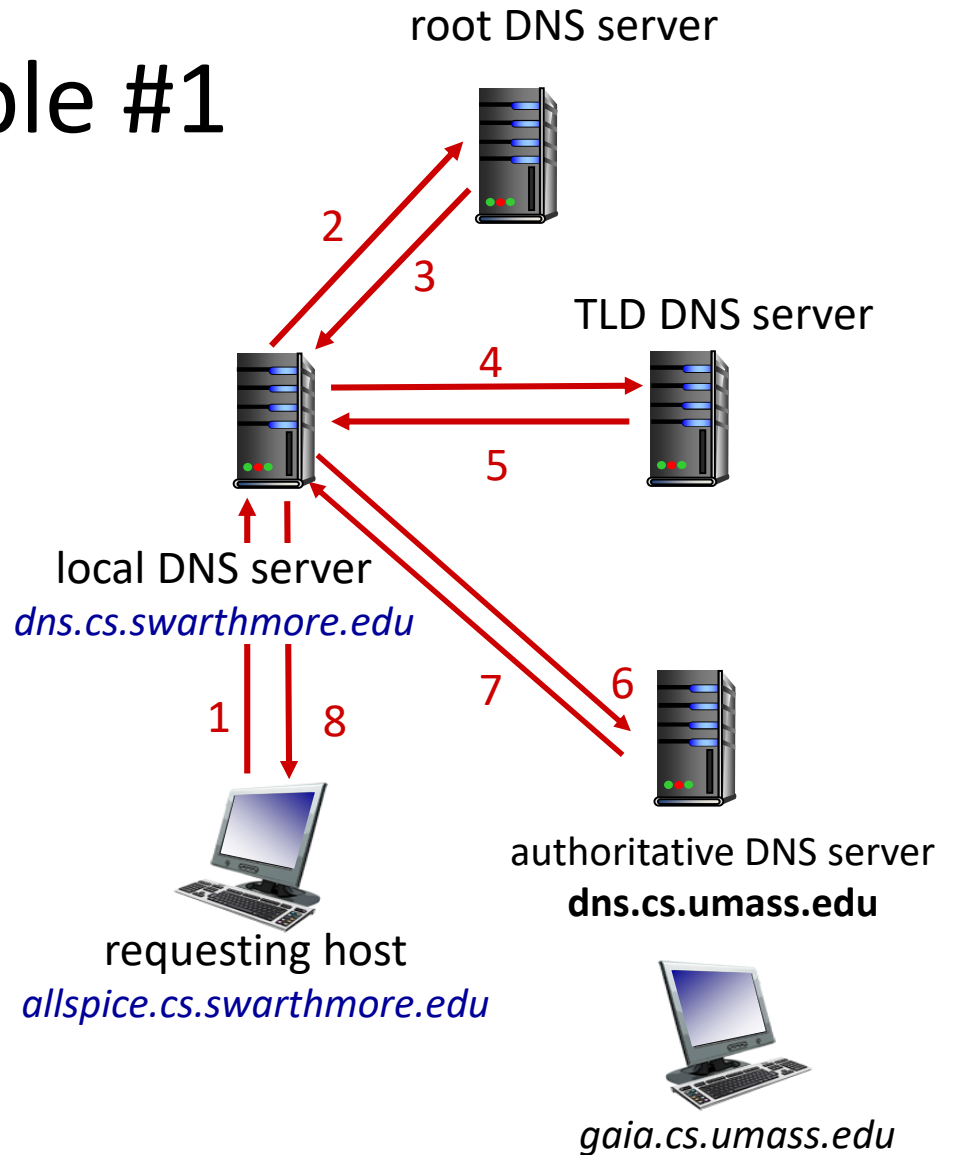


# DNS name resolution example #1

- allspice wants IP address for gaia.cs.umass.edu

## *iterative query:*

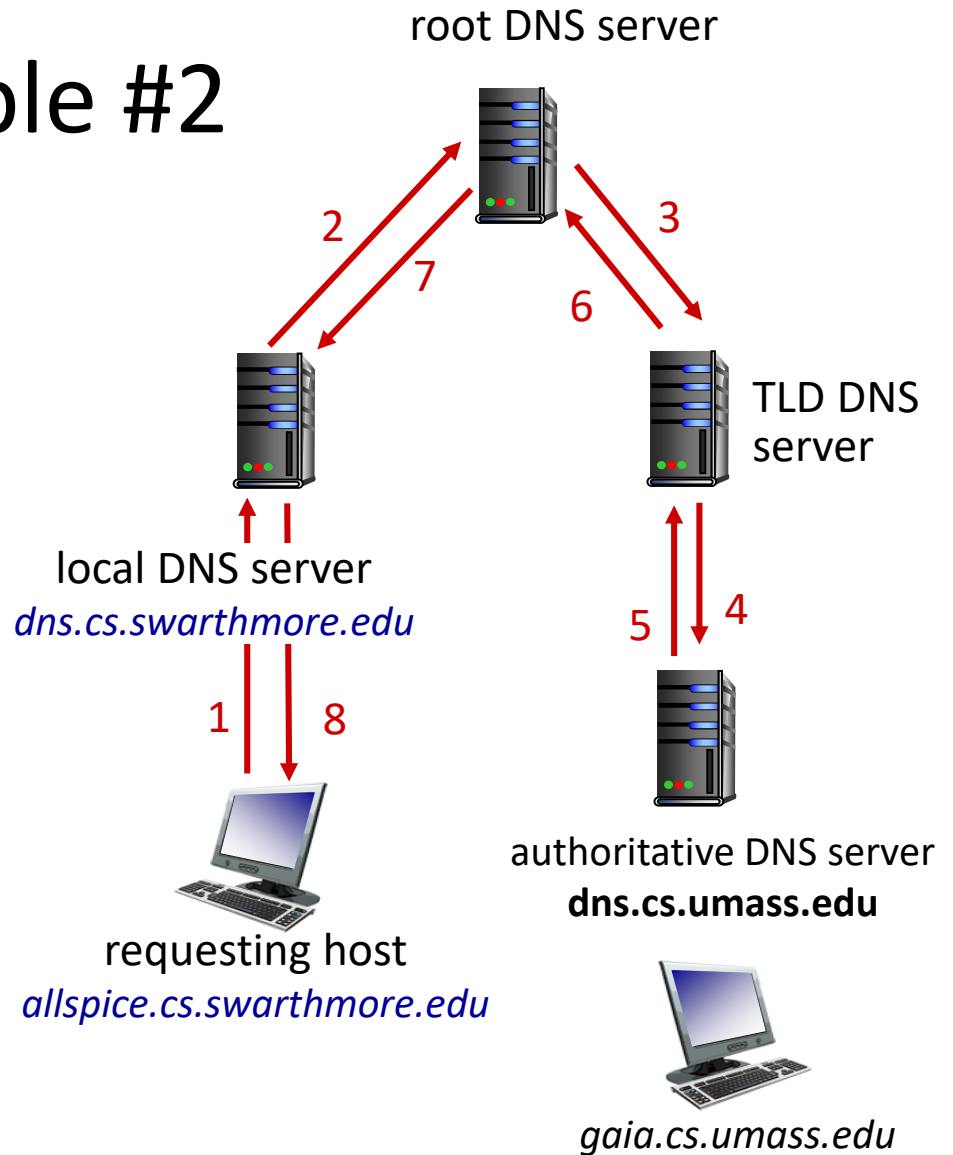
- contacted server replies with name of server to contact
- “I don't know this name, but ask this server”



# DNS name resolution example #2

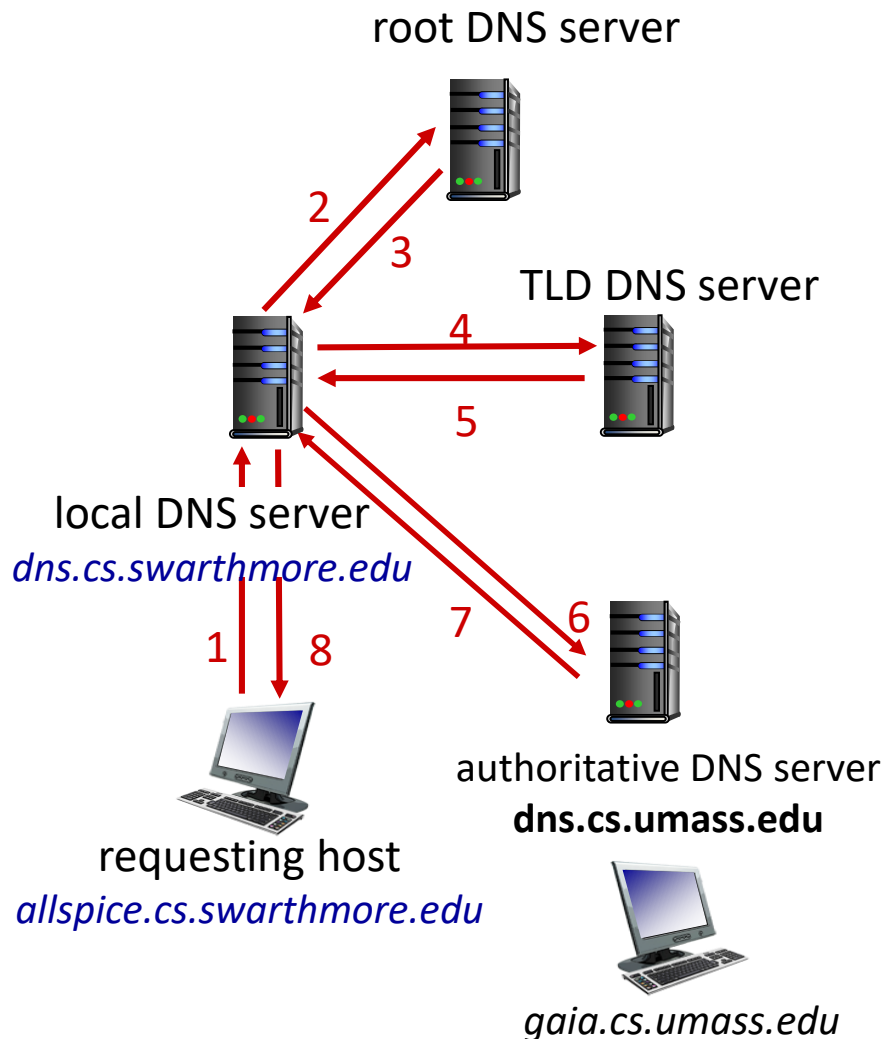
## *recursive query:*

- each server asks the next one, in a chain

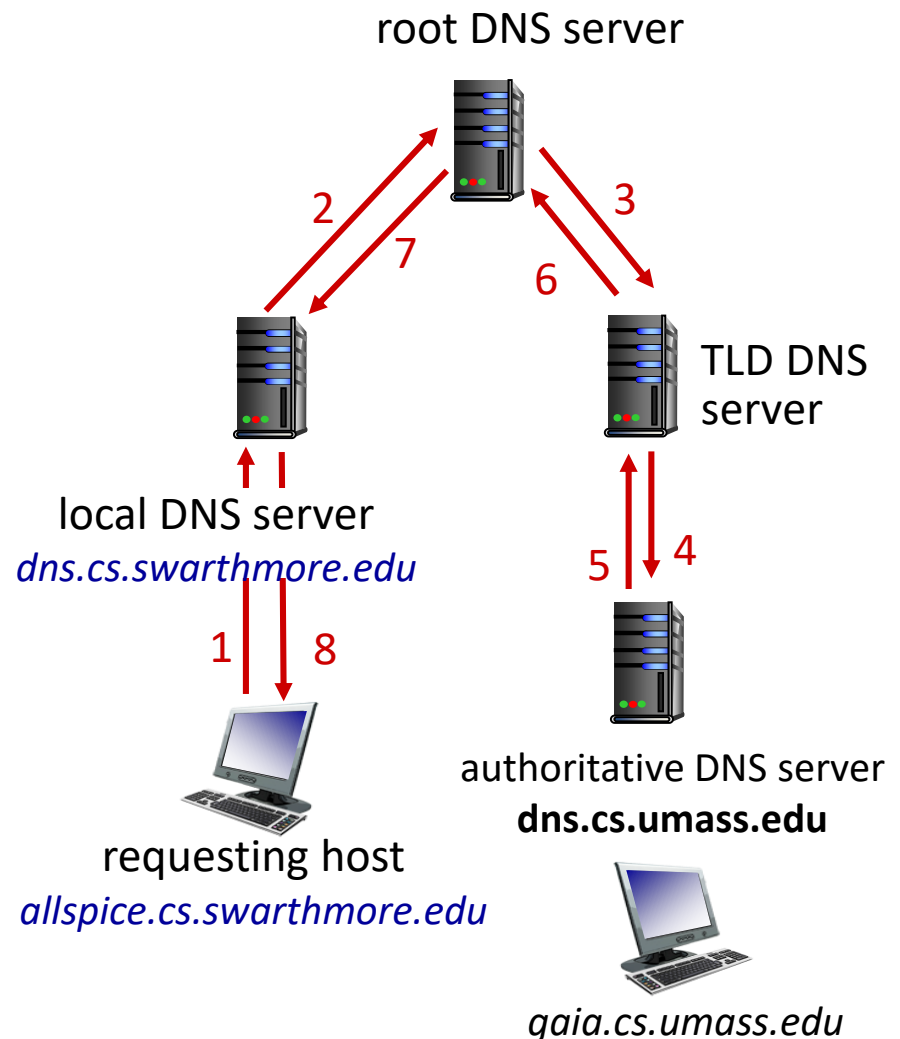


# Which would you use? Why?

## A. Iterative



## B. Recursive



# Caching

- Once (any) name server learns a mapping, it *caches* mapping
  - cache entries timeout (disappear) after some time (TTL: time to live)
  - TLD servers typically cached in local name servers
    - Thus root name servers not often (legitimately) visited

# Caching

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  - TLD servers typically cached in local name servers
    - Thus root name servers not often (legitimately) visited
- (+) Subsequent requests need not burden DNS
- (-) Cached entries may be *out-of-date* (best effort!)
  - If host's name or IP address changes, it may not be known Internet-wide until all TTLs expire

# The TTL value should be

- A. Short, to make sure that changes are accurately reflected
- B. Long, to avoid re-queries of higher-level DNS servers
- C. Something else

# Inserting (or changing) records

- Example: new startup “Network Utopia”
- Register networkuptopia.com at *DNS registrar*
  - provide names, IP addresses of authoritative name server (primary and secondary)
  - registrar inserts two RRs into .com TLD server:  
(networkuptopia.com, dns1.networkuptopia.com, NS)  
(dns1.networkuptopia.com, 212.212.212.1, A)
- Set up authoritative server at that name/address
  - Create records for the services:
    - type A record for www.networkuptopia.com
    - type MX record for @networkuptopia.com email

# DNS Load Balancing

- One load balancing option (others use routing)
- When the authoritative name server responds
  - Round robin between servers
  - Take server load into account?
  - Take location into account (content distribution)



# Summary

- DNS maps human names to IP addresses
- DNS arranged into a hierarchy
  - Scalability / distributed responsibility
  - Autonomous control of local name servers
- Caching crucial for performance